

User's Manual

# GameCube DSP (GDSP)

Reversed and documented by Duddie ([duddie@walla.com](mailto:duddie@walla.com))

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### III. Version history

<b>Version</b>	<b>Date</b>	<b>Author</b>	<b>Change</b>
0.0.1	2005.05.08	Duddie	Initial release
0.0.2	2005.05.09	Duddie	Added \$prod and \$config registers, table of opcodes, disclaimer
0.0.3	2005.05.09	Duddie	Fixed BLOOP and BLOOPI and added description of Loop Stack
0.0.4	2005.05.12	Duddie	Added preliminary DSP memory map and opcode syntax

## **IV. Overview**

## 1. DSP Memory Map

DSP accesses memory in words, so all addresses refer to words. DSP word is 16 bit long.

Instruction Memory (IMEM) is divided into instruction RAM (IRAM) and instruction ROM (IROM).

Exception vectors are located at the top of the RAM and occupy first 8 words.

DSP IRAM is mapped through as first 8kB (4kW) of ARAM (Accelerator RAM) therefore CPU can directly DMA DSP code to DSP IRAM. This usually happens during boottime because DSP ROM is not enabled at cold reset and needs to be reenabled by small stub executed in IRAM.

0x0000	IRAM
0x0fff	
0x8000	IROM
0x8fff	

## V. Registers

## 1. Register names

DSP has 32 16 bit registers although their purpose and their function differ from register to register.

\$0	\$r00	\$ar0	Addressing register 0
\$1	\$r01	\$ar1	
\$2	\$r02	\$ar2	
\$3	\$r03	\$ar3	
\$4	\$r04	\$ix0	
\$5	\$r05	\$ix1	
\$6	\$r06	\$ix2	
\$7	\$r07	\$ix3	
\$8	\$r08		
\$9	\$r09		
\$10	\$r0a		
\$11	\$r0b		
\$12	\$r0c	\$st0	
\$13	\$r0d	\$st1	
\$14	\$r0e	\$st2	
\$15	\$r0f	\$st3	
\$16	\$r10	\$ac0.h	
\$17	\$r11	\$ac1.h	
\$18	\$r12	\$config	
\$19	\$r13	\$sr	
\$20	\$r14	\$prod.l	
\$21	\$r15	\$prod.m1	
\$22	\$r16	\$prod.h	
\$23	\$r17	\$prod.m2	
\$24	\$r18	\$ax0.l	
\$25	\$r19	\$ax1.l	
\$26	\$r1a	\$ax1.h	
\$27	\$r1b	\$ax1.h	
\$28	\$r1c	\$ac0.l	
\$29	\$r1d	\$ac1.l	
\$30	\$r1e	\$ac0.m	
\$31	\$r1f	\$ac1.m	

## 2. Accumulators

DSP has two long 40-bit accumulators (\$acX) and their short 24-bit forms (\$acsX) that reflect upper part of 40-bit accumulator. There are additional two 32-bit accumulators (\$axX).

### Accumulators \$acX:

40-bit accumulator \$acX (\$acX.hml) consists of registers:

$$\$acX = \$acX.h \ll 32 \mid \$acX.m \ll 16 \mid \$acX.l$$

### Short accumulators \$acsX:

24-bit accumulator \$acsX (\$acX.hm) consists of upper 24bit of accumulator \$acX

$$\$acsX = \$acX.h \ll 16 \mid \$acX.m$$

### Additional accumulators \$axX:

$$\$axX = \$axX.h \ll 16 \mid \$axX.l$$

### **3. Stacks**

GDSP contains 4 stack registers:

- \$st0 - call stack
- \$st1 - data stack
- \$st2 - loop address stack
- \$st3 - loop counter

Stacks are implemented in hardware and have limited depth. Data stack is limited to 4 values and call stack is limited to 8 values. Loop stack is limited to 4 values. Upon underflow or overflow of any of the stack registers exception STOVF is raised.

Loop stack is used to control execution of repeated blocks of instructions. Whenever there is value on stack \$st2 and current PC is equal value at \$st2, then value at stack \$st3 is decremented. If value is not zero then PC is modified with value from call stack \$st0. Otherwise values from callstack \$st0 and both loop stacks \$st2 and \$st3 are popped and execution continues at next opcode.

#### ***4. Config register***

It's purpose is unknown at this time. It is written with 0x00ff and 0x0004 values.

## 5. Status register

Status register \$sr reflects flags computed on accumulators after logical or arithmetical operations. Furthermore it also contains control bits to configure flow of certain operations.

Bit	Name	Comment
14	AM	Product multiply result by 2 (when AM = 0)
9	IE	Interrupt enable
8	0	Hardwired to 0 (?)
6	LZ	Logic zero
4	AS	
3	S	Sign
2	Z	Zero

## **6. Product register**

Product register is an intermediate product of multiply or multiply and accumulation. It's result should never be used for calculation although the register can be read or writtent. It reflects state of internal multiply unit. Product is 40 bit with 1 bit of overflow.

```
$prod = ($prod.h << 32) + (($prod.m1 + $prod.m2) << 16) + $prod.l
```

It needs to be noted that \$prod.m1 + \$prod.m2 overflow bit (bit 16) will be added to \$prod.h.

Bit \$sr.AM affects result of multiply unit. If bit \$sr.AM is equal 0 then result of every multiply operation will be multiplied by 2 (two).

## **VI. Exceptions**

## **1. Exception processing**

Exception processing happens by setting program counter to different exception vectors. At the exception time, exception program counter is stored at call stack \$st0 and status register \$sr is stored at data stack \$st1.

### **Operation:**

```
PUSH_STACK($st0)
$st0 = $pc
PUSH_STACK($st1)
$st1 = $sr
$pc = exception_nr * 2
```

## 2. Exception vectors

Exception vectors are located at address 0x0000 in Instruction RAM.

Level	Address	Name	Description
0	0x0000	RESET	
1	0x0002	STOVF	Stack under/overflow
2	0x0004		
3	0x0006		
4	0x0008		
5	0x000a	ACCOV	Accelerator address overflow
6	0x000c		
7	0x000e		

## **VII. Hardware interface**

## 1. Hardware registers

Hardware registers occupy address space at 0xffff in DSP memory space. Each register is 16 bit.

Address	Name	Description
<i>Mailboxes</i>		
0xffffe	CMBH	CPU Mailbox H
0xfffff	CMBL	CPU Mailbox L
0xffffc	DMBH	DSP Mailbox H
0xffffd	DMBL	DSP Mailbox L
<i>DMA interface</i>		
0xffce	DSMAH	Memory address H
0xffcf	DSMAL	Memory address L
0xffcd	DSPA	DSP memory address
0xffc9	DSCR	DMA Control
0xffcb	DSBL	Block size
<i>Accelerator</i>		
0xffd4	ACSAH	Accelerator start address H
0xffd5	ACSAL	Accelerator start address L
0xffd6	ACEAH	Accelerator end address H
0xffd7	ACEAL	Accelerator end address L
0xffd8	ACCAH	Accelerator current address H
0xffd9	ACCAL	Accelerator current address L
0xffdd	ACDAT	Accelerator data
<i>Interrupts</i>		
0xffffb	DIRQ	IRQ request

## 2. Interrupts

DSP can raise interrupts at CPU. Usually interrupts are used to signal that new DSP mbox has been filled with data.

<b>0xFFFB</b>	<b>DIRQ</b>	<b>IRQ Request</b>
-----I		

Bit	Name	R/W	Action
0	I	W	1 - Raise interrupt at CPU

### 3. Mailboxes

CPU Mailbox (CMB) is a register that allows sending 31 bits of information from CPU to DSP.

<b>0xFFFFE</b>	<b>CMBH</b>	<b>CPU Mailbox H</b>
Mddd dddd dddd dddd		

Bit	Name	R/W	Action
15	M	R	1 - Mailbox contains mail from CPU 0 - Mailbox empty
14-0	d	R	bits 30-16 of mail from CPU

<b>0xFFFF</b>	<b>CMBL</b>	<b>CPU Mailbox L</b>
dddd dddd dddd dddd		

Bit	Name	R/W	Action
15-0	d	R	bits 15-0 of mail from CPU. Reading this register by DSP causes M bit of register CMBH to be cleared.

#### Operation:

From CPU side, software usually checks M bit of CMBH. It takes action only in case this bit is 0. Action is to write CMBH first and then CMBL. After writing CMBL mail is ready to be received by DSP.

From DSP side, DSP loops by probing M bit. When this bit is 1 it reads CMBH first and then CMBL. After reading CMBL bit M of CMBH signaling mail from CPU will be cleared.

DSP mailbox (DMB) is an interface to send 31 bits of information from DSP to CPU.

<b>0xFFFC</b>	<b>DMBH</b>	<b>DSP Mailbox H</b>
Mddd dddd dddd dddd		

Bit	Name	R/W	Action
15	M	R	1 – Mailbox has not been received by CPU 0 – Mailbox empty
		W	Does not matter. It will be set when DMBL is written.
14-0	d	W	bits 30-16 of mail from DSP to CPU

<b>0xFFFD</b>	<b>DMBL</b>	<b>DSP Mailbox L</b>
dddd dddd dddd dddd		

Bit	Name	R/W	Action
15-0	d	W	bits 15-0 of mail from DSP to CPU. Writing this register by DSP causes M bit of register DMBH to be set signaling that mail is ready.

### Operation:

Sending mail from DSP to CPU can be achieved by writing mail to DMBH and then to DMBL registers. After writing DMBL a flag M in DMBH will be set signalling that mail is ready to be received by CPU. If DSP needs to receive response from CPU then it usually waits for bit M to be cleared after sending a mail. If DSP does processing when CPU receives a mail, then it waits for bit M to be cleared before issuing another mail to CPU.

## 4. DMA

GDSP is connected with memory bus through DMA channel. DMA can transfer data between DSP memory (both instruction and data) and main memory.

<b>0xFFCE</b>	<b>DSMAH</b>	<b>Memory Address H</b>
dddd dddd dddd dddd		

Bit	Name	R/W	Action
15-0	d	R	bits 31-16 of main memory address

<b>0xFFCF</b>	<b>DSMAL</b>	<b>Memory address L</b>
dddd dddd dddd dddd		

Bit	Name	R/W	Action
15-0	d	R	bits 15-0 of main memory address

<b>0xFFCD</b>	<b>DSPA</b>	<b>DSP Address</b>
dddd dddd dddd dddd		

Bit	Name	R/W	Action
15-0	d	W	bits 15-0 of DSP memory address

<b>0xFFCB</b>	<b>DSBL</b>	<b>DSP Address</b>
dddd dddd dddd dddd		

Bit	Name	R/W	Action
15-0	d	W	length in bytes of transfer. writing to this register starts DMA transfer.

<b>0xFFC9</b>	<b>DSCR</b>	<b>DSP Address</b>
-----		

<b>Bit</b>	<b>Name</b>	<b>R/W</b>	<b>Action</b>
15-0	d	W	

## **5. Accelerator**

Accelerator is used to transfer data from accelerator memory (ARAM) to DSP. Accelerator area can be marked with ACSA (start) and ACEA (end) addresses. Current address for can be set or read from ACCA register. Reading from accelerator memory is done by reading from ACDAT register. This register contains data from ARAM pointed by ACCA register. After reading, ACCA is incremented by one. After ACCA grows bigger than area pointed by ACEA, it gets reset to a value from ACSA and ACCOV interrupt is generated.

## VIII. Opcodes

## 1. Opcode syntax

### Basic syntax of opcode:

OPC                   opc\_params

*Above syntax is correct for all opcodes.*

OPC                   - opcode  
opc\_params           - opcode parameters if necessary

### EXAMPLES:

```
JMP           0x0300  
CALL          loop  
HALT
```

### Extended syntax:

OPC'EXOPC   opc\_params : exopc\_params

*Above syntax is correct only for arithmetic opcodes because those can be extended with additional load/store unit behaviour.*

OPC                   - opcode  
OPC                   - extended opcode  
opc\_params           - opcode parameters if necessary  
opc\_params           - opcode parameters for extended part if  
                          necessary

### EXAMPLES:

```
DECM'L       $acs0 : $ac1.m, @ar0  
NX'MV        : $acx1.h, $ac0.l
```

## 2. Operation - used functions

Functions used for describing operation of opcodes

`PUSH_STACK($stR)`

### **Description:**

Pushes value onto given stack referenced by stack register `$stR`. Operation moves down values in internal stack.

### **Operation:**

```
stack_stR[stack_ptr_stR++] = $stR;
```

`POP_STACK($stR)`

### **Description:**

Pops value from stack referenced by stack register `$stR`. Operation moves values up in internal stack.

### **Operation:**

```
$stR = stack_stR[--stack_ptr_stR]
```

`FLAGS(val)`

### **Description:**

Calculates flags depending on given value or result of operation and setting corresponding bits in status register `$sr`.

### **Operation:**

EXECUTE\_OPCODE(new\_pc)

**Description:**

Executes opcode at given new\_pc address.

**Operation:**

### 3. Meaning of bits

Opcode decoding uses special naming for bits and their decimal representations to provide easier understanding of bit fields in opcode

Binary form	Decimal form	Meaning
d, dd, ddd, dddd	D	Destination register
s, ss, sss, ssss	S	Source register
t, tt, ttt, tttt	T	Source register
r, rr, rrr, rrrr	R	Register (either source or destination)
Aaaaa(a)	A, addrA	Address in either I or D memory
xxxx xxxx	X	Extended opcode
mmm(m)	M, addrM	Address in memory
iii(i)	I, Imm	Immediate value
cccc	cc	Condition (See conditional opcodes)

#### 4. Conditional opcodes

Conditional opcodes are being executed only when given condition described by conditional field has been met. To the group of conditional opcodes belong: CALL, JMP, IF, RET.

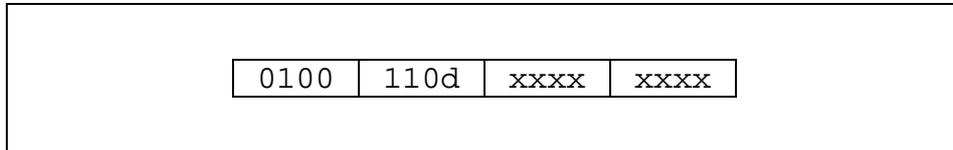
Bits	cc	Name	Evaluated expression
0000			
0001			
0010			
0011			
0100	EQ	Equal	
0101	NE	Not equal	
0110			
0111			
1000			
1001			
1010			
1011			
1100	ZR	Zero	$\$sr \& 0x40$
1101	NZ	Not zero	$!(\$sr \& 0x40)$
1110			
1111		<always>	

**Note:**

There is two pairs of conditions that work similar: EQ/NE and ZR/NZ. EQ/NE pair operates on arithmetic zero flag (arithmetic 0) while ZR/NZ pair operates on logic zero flag (logic 0).

## ***5. Opcodes decoding***

# ADD



## Format:

ADD        \$acD, \$ac(1-D)

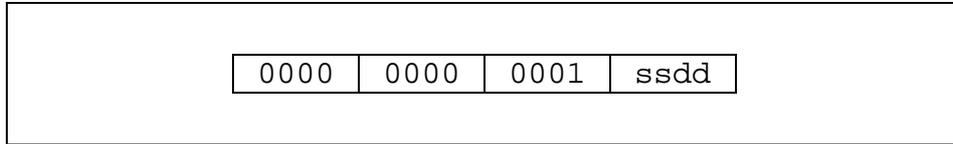
## Description:

Adds accumulator \$ac(1-D) to accumulator register \$acD.

## Operation:

```
$acD += $ac(1-D)
FLAGS($acD)
$pc++
```

# ADDARN



## Format:

ADDARN \$arD, \$ixS

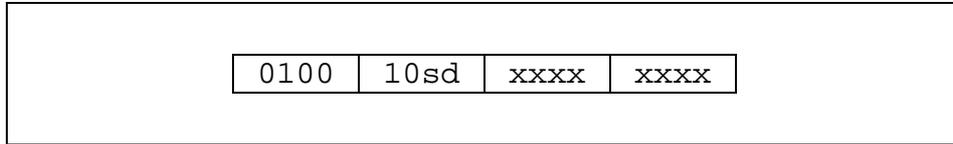
## Description:

Adds indexing register \$ixS to an addressing register \$arD.

## Operation:

```
$arD += $ixS  
$pc++
```

# ADDAX



## Format:

ADDAX     \$acD, \$axS

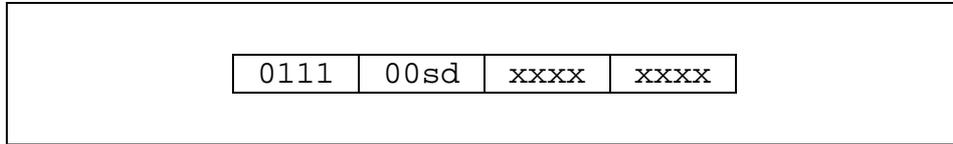
## Description:

Adds secondary accumulator \$axS to accumulator register \$acD.

## Operation:

```
$acD += $axS  
FLAGS($acD)  
$pc++
```

# ADDAXL



## Format:

ADDAXL    \$acD, \$axS.l

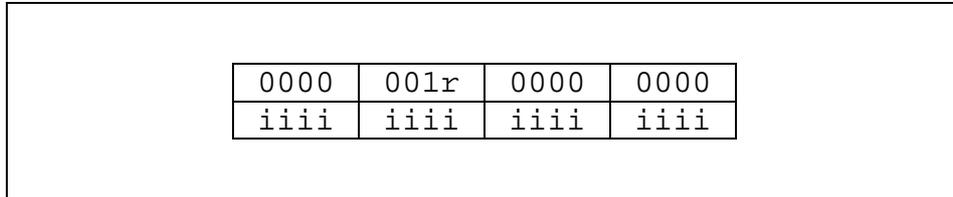
## Description:

Adds secondary accumulator \$axS.l to accumulator register \$acD.

## Operation:

```
$acD += $axS.l  
FLAGS($acD)  
$pc++
```

# ADDI



## Format:

ADDI      \$amR, #I

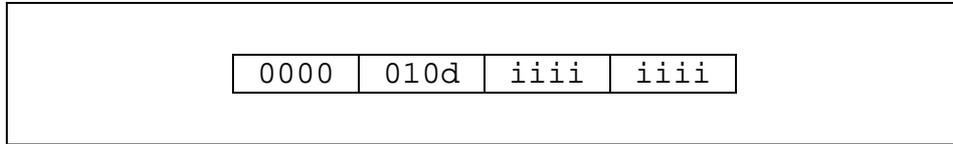
## Description:

Adds immediate (16-bit sign extended) to mid accumulator \$acD.hm.

## Operation:

```
$acD.hm += #I  
FLAGS($acD)  
$pc++
```

# ADDIS



## Format:

ADDIS     \$acD, #I

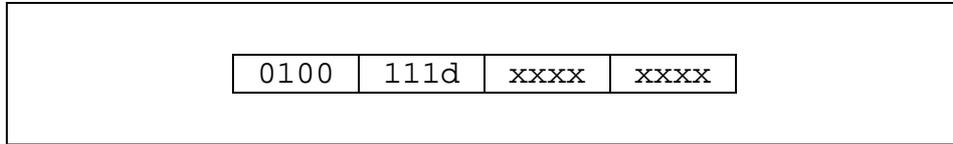
## Description:

Adds short immediate (8-bit sign extended) to mid accumulator \$acD.hm.

## Operation:

```
$acD.hm += #I  
FLAGS($acD)  
$pc++
```

# ADDP



## Format:

ADDP      \$acD

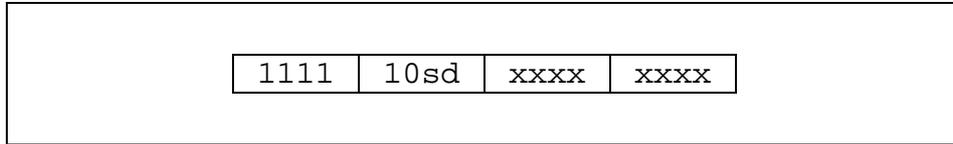
## Description:

Adds product register to accumulator register.

## Operation:

```
$acD += $prod  
FLAGS($acD)  
$pc++
```

# ADDPAXZ



## Format:

```
ADDPAXZ $acD, $axS
```

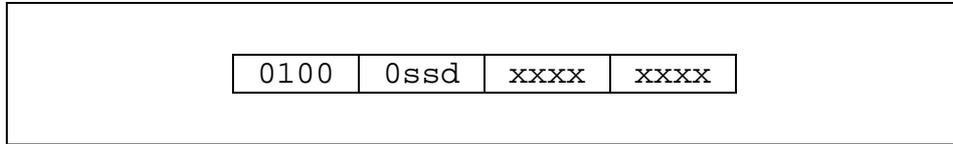
## Description:

Adds secondary accumulator \$axS to product register and stores result in accumulator register. Low 16-bits of \$acD (\$acD.l) are set to 0.

## Operation:

```
$acD.hm = $prod.hm + $ax.h  
$acD.l = 0  
FLAGS($acD)  
$pc++
```

# ADDR



## Format:

ADDR      \$acD, \$(0x18+S)

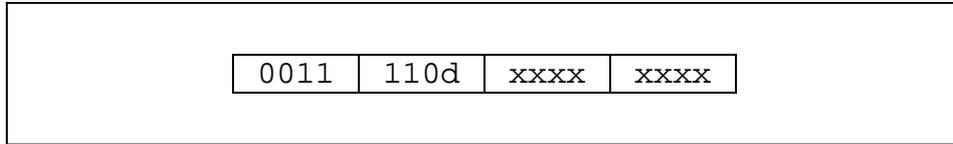
## Description:

Adds register \$(0x18+S) to accumulator \$acD register.

## Operation:

```
$acD += $(0x18+S)
FLAGS($acD)
$pc++
```

# ANDC



## Format:

AMDC      \$acD.m, \$ac(1-D).m

## Description:

Logic AND middle part of accumulator \$acD.m with middle part of accumulator \$ax(1-D).m.

## Operation:

```
$acD.m &= $ac(1-D).m  
FLAGS($acD)  
$pc++
```

# ANDCF

0000	001r	1010	0000
iiii	iiii	iiii	iiii

## Format:

ANDCF     \$acD.m, #I

## Description:

Set logic zero (LZ) flag in status register \$sr if result of logical AND operation of accumulator mid part \$acD.m with immediate value I is equal immediate value I.

## Operation:

```
IF    ($acD.m & I) == I
      $sr.LZ = 1
ELSE
      $sr.LZ = 0
$pc++
```

# ANDF

0000	001r	1100	0000
iiii	iiii	iiii	iiii

## Format:

ANDF      \$acD.m, #I

## Description:

Set logic zero (LZ) flag in status register \$sr if result of logic AND of accumulator mid part \$acD.m with immediate value I is equal zero.

## Operation:

```
IF    ($acD.m & I) == 0
      $sr.LZ = 1
ELSE
      $sr.LZ = 0
$pc++
```

# ANDI

0000	001r	0100	0000
iiii	iiii	iiii	iiii

## Format:

ANDI        \$acD.m, #I

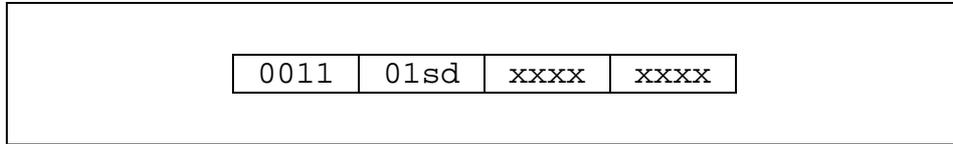
## Description:

Logic AND of accumulator mid part \$acD.m with immediate value I.

## Operation:

```
$acD.m &= #I  
FLAGS($acD)  
$pc++
```

# ANDR



## Format:

ANDR      \$acD.m, \$axS.h

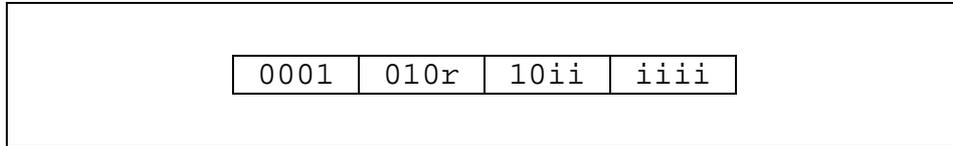
## Description:

Logic AND middle part of accumulator \$acD.m with high part of secondary accumulator \$axS.h.

## Operation:

```
$acD.m &= $axS.h  
FLAGS($acD)  
$pc++
```

# ASL



## Format:

ASL        \$acR, #I

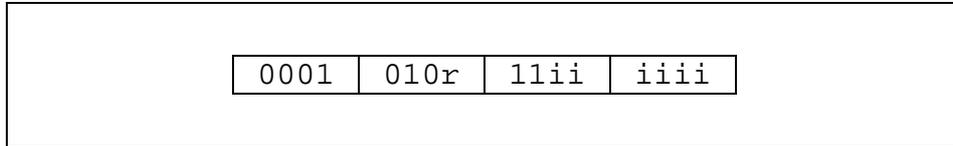
## Description:

Logically shifts left accumulator \$acR by number specified by value I.

## Operation:

\$acR <<= I  
FLAGS(\$acD)  
\$pc++

# ASR



## Format:

ASR      \$acR, #I

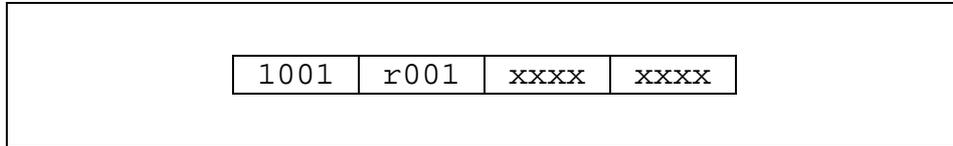
## Description:

Arithmetically shifts left accumulator \$acR by number specified by value calculated by negating sign extended bits 0-6.

## Operation:

```
$acR <<= I  
FLAGS($acD)  
$pc++
```

# ASR16



## Format:

ASR16      \$acR

## Description:

Arithmetically shifts right accumulator \$acR by 16.

## Operation:

```
$acR >>= 16  
FLAGS($acD)  
$pc++
```

# BLOOP

0000	0000	011r	rrrr
aaaa	aaaa	aaaa	aaaa

## Format:

BLOOP     \$R, addrA

## Description:

Repeatedly execute block of code starting at following opcode until counter specified by value from register \$R reaches zero. Block ends at specified address addrA inclusive, ie. opcode at addrA is the last opcode included in loop. Counter is pushed on loop stack \$st3, end of block address is pushed on loop stack \$st2 and repeat address is pushed on call stack \$st0. Up to 4 nested loops is allowed.

## Operation:

```
$st0 = $pc + 2
$st2 = addrA
$st3 = $R
$pc + 2
// in real hardware below does not happen, this
opcode only sets stack registers
WHILE ($st3-- )
  DO
    EXECUTE_OPCODE($pc)
    WHILE($pc != $st2)
      $pc = $st0
  $pc = addrA + 1
// remove vaues from stack
```

## See also:

Description of Stack registers explains how loop stacks are working

# BLOOPI

0001	0001	iiii	iiii
aaaa	aaaa	aaaa	aaaa

## Format:

BLOOPI #I, addrA

## Description:

Repeatedly execute block of code starting at following opcode until counter specified by immediate value I reaches zero. Block ends at specified address addrA inclusive, ie. opcode at addrA is the last opcode included in loop. Counter is pushed on loop stack \$st3, end of block address is pushed on loop stack \$st2 and repeat address is pushed on call stack \$st0. Up to 4 nested loops is allowed.

## Operation:

```
$st0 = $pc + 2
$st2 = addrA
$st3 = I
$pc + 2
// in real hardware below does not happen, this
opcode only sets stack registers
WHILE ($st3--)
    DO
        EXECUTE_OPCODE($pc)
        WHILE($pc != $st2)
            $pc = $st0
        $pc = addrA + 1
    // remove vaues from stack
```

## See also:

Description of Stack registers explains how loop stacks are working

# CALL

0000	0010	1011	1111
aaaa	aaaa	aaaa	aaaa

## Format:

CALL        addressA

## Description:

Call function. Push program counter of instruction following "call" to call stack \$st0. Set program counter to address represented by value that follows this "call" instruction.

## Operation:

```
// must skip value that follows "call"  
PUSH_STACK($st0)  
$st0 = $pc + 2  
$pc = addressA
```

# CALLcc

0000	0010	1011	cccc
aaaa	aaaa	aaaa	aaaa

## Format:

CALLcc      addressA

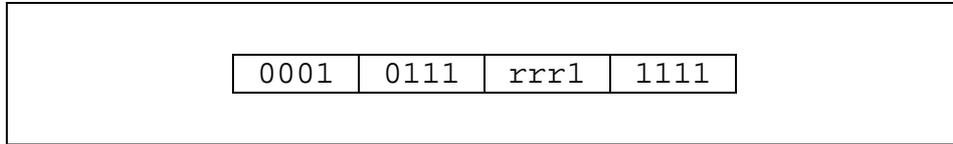
## Description:

Call function if condition cc has been met. Push program counter of instruction following "call" to call stack \$st0. Set program counter to address represented by value that follows this "call" instruction.

## Operation:

```
// must skip value that follows "call"
IF (cc)    PUSH_STACK($st0)
           $st0 = $pc + 2
           $pc = addressA
ELSE
           $pc += 2
```

# CALLR



## Format:

CALLR     \$R

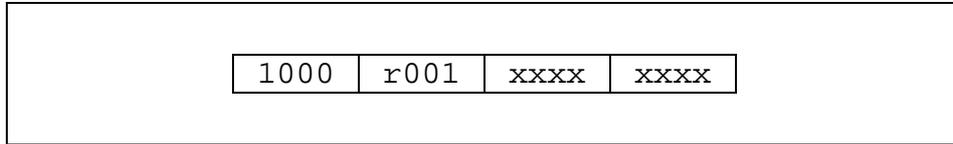
## Description:

Call function. Push program counter of instruction following "call" to call stack \$st0. Set program counter to register \$R.

## Operation:

```
PUSH_STACK($st0)
$st0 = $pc + 1
$pc = $R
```

# CLR



## Format:

CLR      \$acR

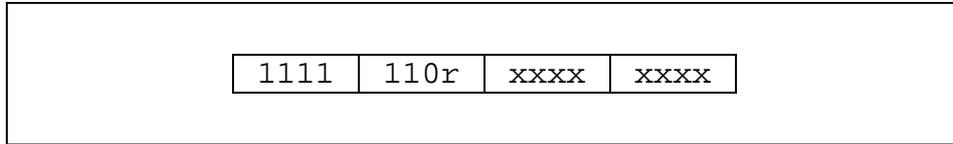
## Description:

Clears accumulator \$acR

## Operation:

```
$acR = 0  
FLAGS($acR)  
$pc++
```

# CLRL



## Format:

CLRD      \$acR.l

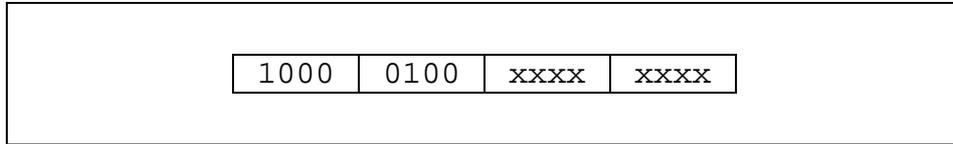
## Description:

Clears \$acR.l - low 16 bits of accumulator \$acR.

## Operation:

\$acR.l = 0  
FLAGS(\$acR)  
\$pc++

# CLRP



## Format:

CLRP

## Description:

Clears product register \$prod.

## Operation:

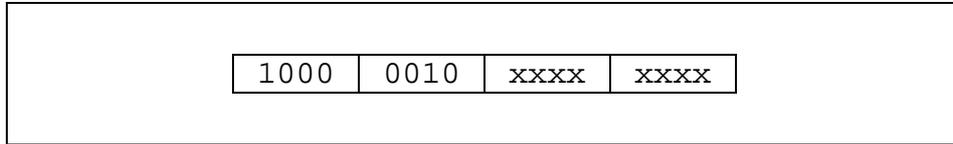
```
$prod = 0 // see note below  
$pc++
```

## Note:

Actually product register gets cleared by setting registers with following values:

```
$14 = 0x0000  
$15 = 0xffff0  
$16 = 0x00ff  
$17 = 0x0010
```

# CMP



## Format:

CMP

## Description:

Compares accumulator \$ac0 with accumulator \$ac1.

## Operation:

```
$sr = FLAGS($ac0 - $ac1)  
$pc++
```

# CMPI

0000	001r	1000	0000
iiii	iiii	iiii	iiii

## Format:

CMPI      \$amD, #I

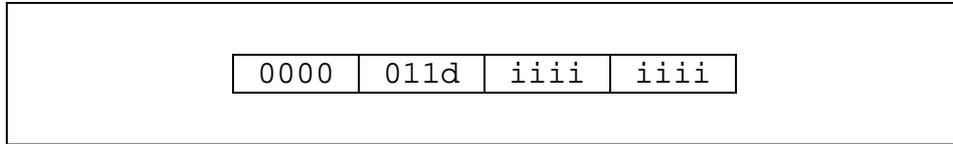
## Description:

Compares mid accumulator \$acD.hm (\$amD) with sign extended immediate value I. Although flags are being set regarding whole accumulator register.

## Operation:

```
res = ($acD.hm - I) | $acD.l  
FLAGS(res)  
$pc++
```

# CMPIS



## Format:

CMPIS      \$acD, #I

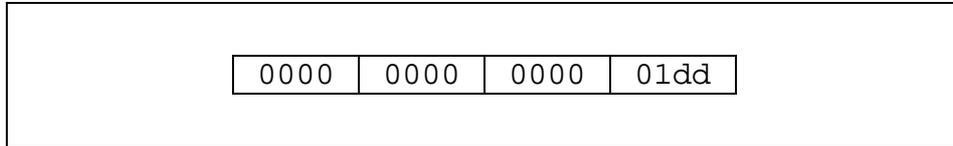
## Description:

Compares accumulator with short immediate. Comparison is executed by subtracting short immediate (8bit sign extended) from mid accumulator \$acD.hm and computing flags based on whole accumulator \$acD.

## Operation:

FLAGS(\$acD - #I)  
\$PC++

# DAR



## Format:

DAR      \$arD

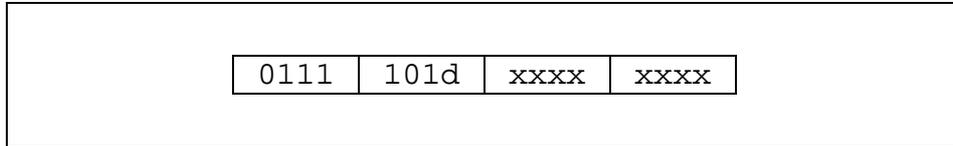
## Description:

Decrement address register \$arD.

## Operation:

\$arD--  
\$pc++

# DEC



## Format:

DEC      \$acD

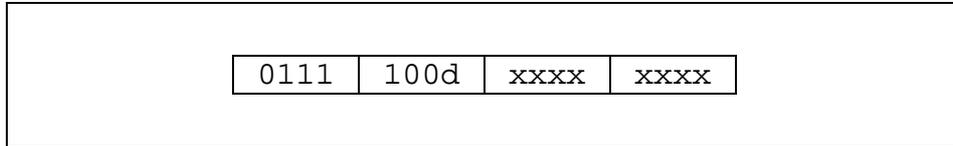
## Description:

Decrement accumulator \$acD.

## Operation:

```
$acD--;  
FLAGS($acD);  
$pc++;
```

# DECM



## Format:

DECM      \$acsD

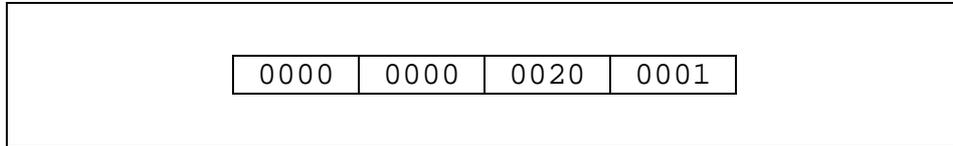
## Description:

Decrement 24-bit mid-accumulator \$acsD.

## Operation:

```
$acsD-- ;  
FLAGS($acD) ;  
$pc++ ;
```

# HALT



## Format:

HALT

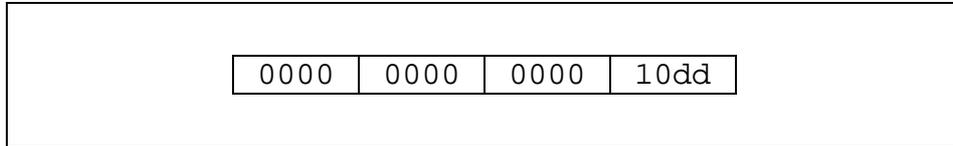
## Description:

Stops execution of DSP code. Sets bit DSP\_CR\_HALT in register DREG\_CR.

## Operation:

$DREG\_CR \mid = DSP\_CR\_HALT;$

# IAR



## Format:

IAR      \$arD

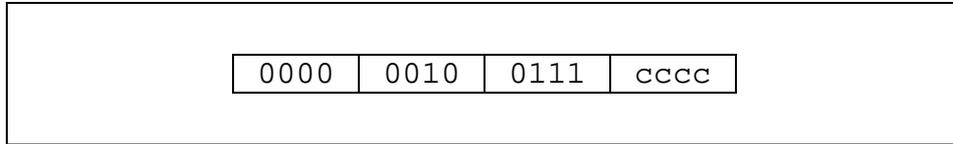
## Description:

Increment address register \$arD.

## Operation:

\$arD++  
\$pc++

# IFcc



## Format:

IFcc

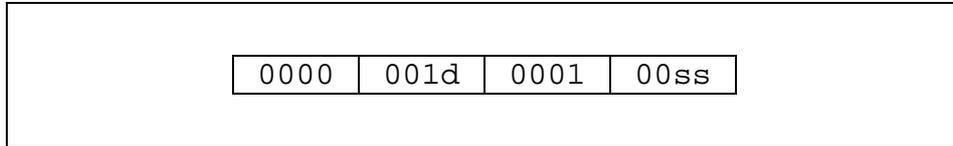
## Description:

Execute following opcode if the condition has been met.

## Operation:

```
IF (cc) EXECUTE_OPCODE($pc + 1)
ELSE    $pc += 2
```

# ILRR



## Format:

ILRR      \$acD.m, @\$arS

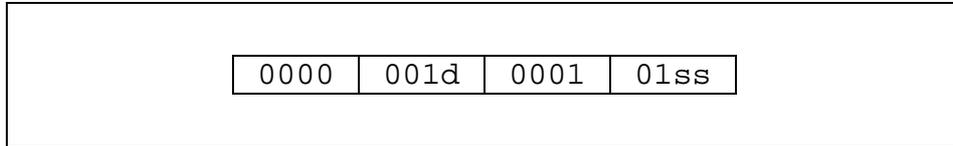
## Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m.

## Operation:

```
$acD.m = MEM[$arS]  
$pc++
```

# ILRRD



## Format:

ILRRD      \$acD.m, @\$arS

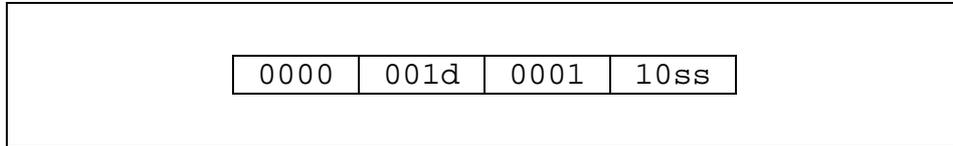
## Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m. Decrement addressing register \$arS.

## Operation:

```
$acD.m = MEM[$arS]  
$arS--  
$pc++
```

# ILRRI



## Format:

ILRRI \$acD.m, @\$S

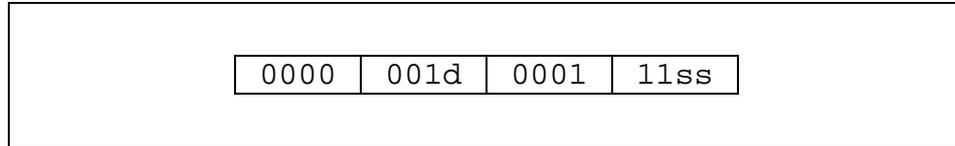
## Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m. Increment addressing register \$arS.

## Operation:

```
$acD.m = MEM[$arS]  
$arS++  
$pc++
```

# ILRRN



## Format:

ILRRN      \$acD.m, @\$arS

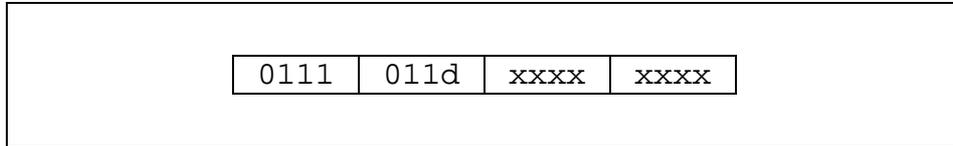
## Description:

Move value from instruction memory pointed by addressing register \$arS to mid accumulator register \$acD.m. Add corresponding indexing register \$ixS to addressing register \$arS.

## Operation:

```
$acD.m = MEM[$arS]  
$arS += $ixS  
$pc++
```

# INC



## Format:

INC      \$acD

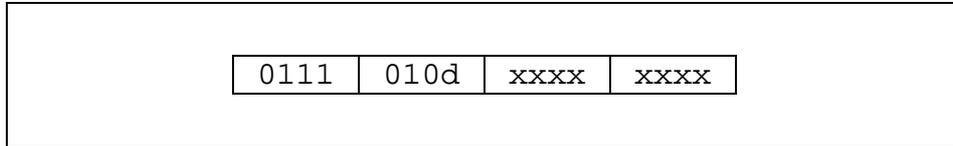
## Description:

Increment accumulator \$acD.

## Operation:

\$acD++  
FLAGS(\$acD)  
\$pc++

# INCM



## Format:

INCM      \$acsD

## Description:

Increment 24-bit mid-accumulator \$acsD.

## Operation:

\$acsD++  
FLAGS (\$acD)  
\$pc++

# JMP

0000	0010	1001	1111
aaaa	aaaa	aaaa	aaaa

## Format:

JMP          addressA

## Description:

Jump to addressA. Set program counter to address represented by value that follows this "jmp" instruction.

## Operation:

\$pc = addressA

# Jcc

0000	0010	1001	cccc
aaaa	aaaa	aaaa	aaaa

## Format:

Jcc            addressA

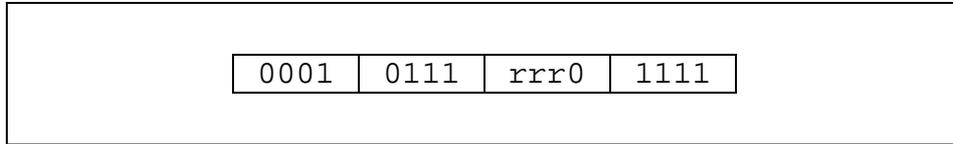
## Description:

Jump to addressA if condition cc has been met. Set program counter to address represented by value that follows this "jmp" instruction.

## Operation:

```
IF (cc)        $pc = addressA
ELSE            $pc += 2
```

# JMPR



## Format:

JMP      \$R

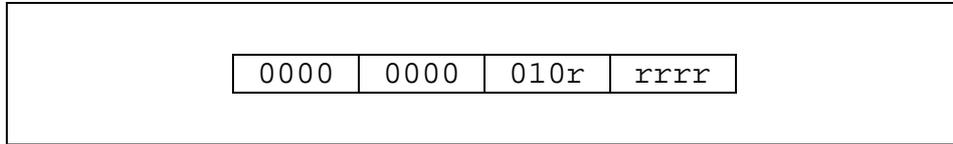
## Description:

Jump to address; set program counter to a value from register \$R.

## Operation:

$\$pc = \$R$

# LOOP



## Format:

LOOP      \$R

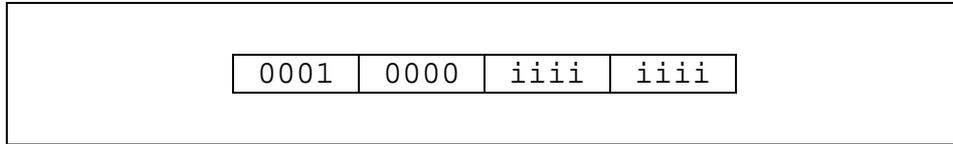
## Description:

Repeatedly execute following opcode until counter specified by value from register \$R reaches zero. Each execution decrement counter. Register \$R remains unchanged. If register \$R is set to zero at the beginning of loop then looped instruction will not get executed.

## Operation:

```
counter = $R
WHILE (counter-- )
    EXECUTE_OPCODE($pc+1)
$pc += 2
```

# LOOPI



## Format:

```
LOOPI    #I
```

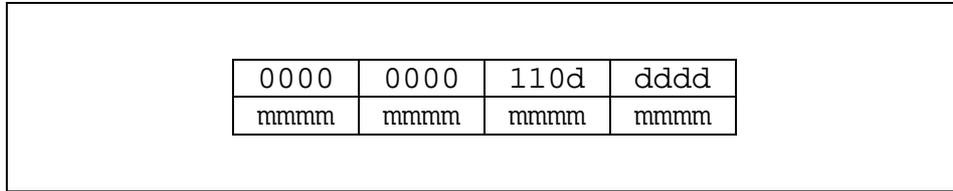
## Description:

Repeatedly execute following opcode until counter specified by immediate value I reaches zero. Each execution decrement counter. If immediate value I is set to zero at the beginning of loop then looped instruction will not get executed.

## Operation:

```
counter = I  
WHILE (counter--)  
    EXECUTE_OPCODE($pc+1)  
$pc += 2
```

# LR



## Format:

LR            \$D, @M

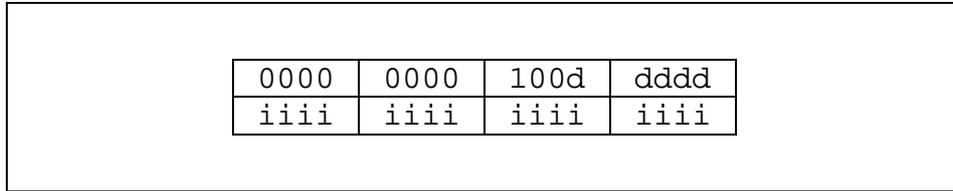
## Description:

Move value from data memory pointed by address M to register \$D.  
Perform additional operation depending on destination register.

## Operation:

\$D = MEM[M]  
\$pc += 2

# LRI



## Format:

LRI            \$D, #I

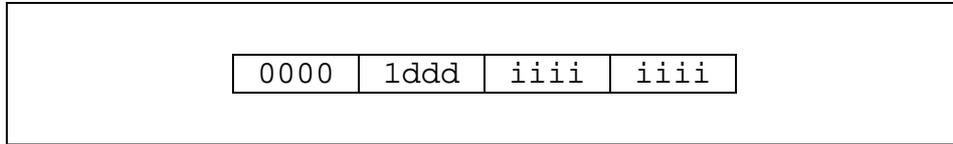
## Description:

Load immediate value I to register \$D. Perform additional operation depending on destination register.

## Operation:

\$D = I  
\$pc += 2

# LRIS



## Format:

LRIS      \$(0x18+D), #I

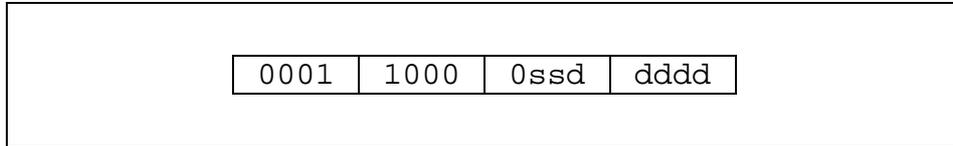
## Description:

Load immediate value I (8-bit sign extended) to accumulator register \$(0x18+D). Perform additional operation depending on destination register.

## Operation:

\$(0x18+D) = I  
\$pc++

# LRR



## Format:

LRR      \$D, @\$S

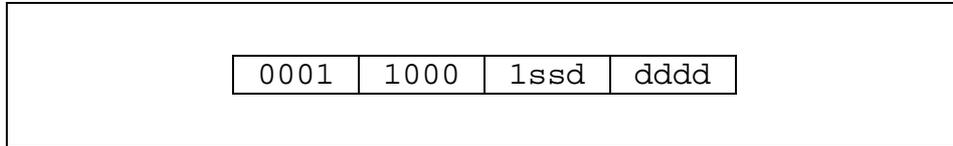
## Description:

Move value from data memory pointed by addressing register \$S to register \$D. Perform additional operation depending on destination register.

## Operation:

\$D = MEM[\$S]  
\$pc++

# LRRD



## Format:

LRRD      \$D, @\$S

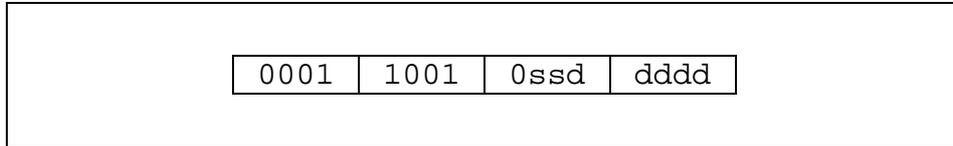
## Description:

Move value from data memory pointed by addressing register \$S to register \$D. Decrement register \$S. Perform additional operation depending on destination register.

## Operation:

\$D = MEM[\$S]  
\$S--  
\$pc++

# LRRI



## Format:

LRRI      \$D, @\$S

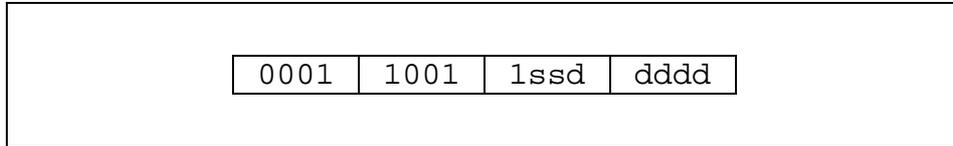
## Description:

Move value from data memory pointed by addressing register \$S to register \$D. Increment register \$S. Perform additional operation depending on destination register.

## Operation:

\$D = MEM[\$S]  
\$S++  
\$pc++

# LRRN



## Format:

LRRN      \$D, @\$S

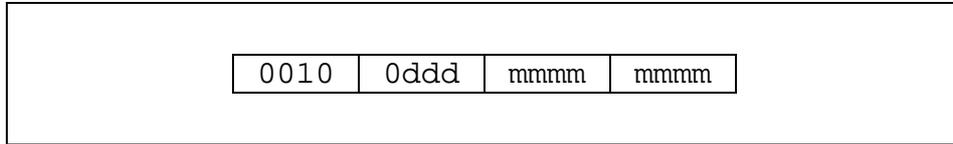
## Description:

Move value from data memory pointed by addressing register \$S to register \$D. Add indexing register \$(0x4+S) to register \$S. Perform additional operation depending on destination register.

## Operation:

```
$D = MEM[$S]  
$S += $(4+S)  
$pc++
```

# LRS



## Format:

LRS      \$(0x18+D), @M

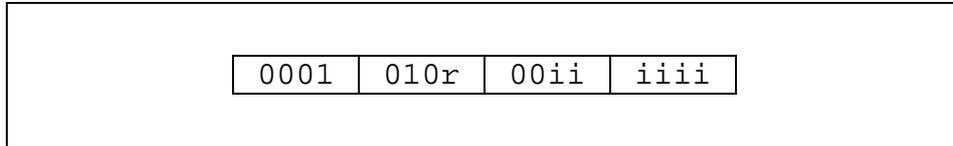
## Description:

Move value from data memory pointed by address M (8-bit sign extended) to register \$(0x18+D). Perform additional operation depending on destination register.

## Operation:

$\$(0x18+D) = \text{MEM}[M]$   
 $\$pc += 2$

# LSL



## Format:

LSL          \$acR, #I

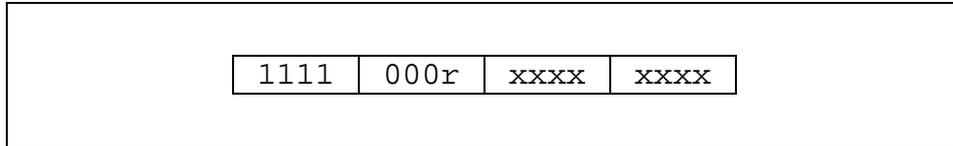
## Description:

Logically shifts left accumulator \$acR by number specified by value I.

## Operation:

\$acR <<= I  
FLAGS(\$acD)  
\$pc++

# LSL16



## Format:

LSL16\$acR

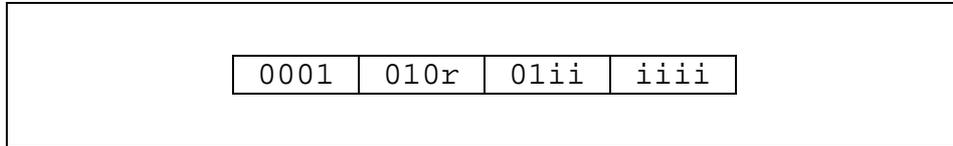
## Description:

Logically shifts left accumulator \$acR by 16.

## Operation:

\$acR <<= 16  
FLAGS(\$acD)  
\$pc++

# LSR



## Format:

LSR      \$acR, #I

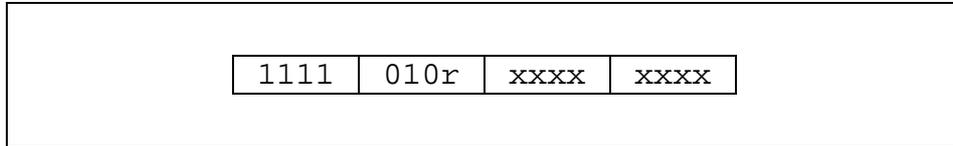
## Description:

Logically shifts left accumulator \$acR by number specified by value calculated by negating sign extended bits 0-6.

## Operation:

```
$acR <<= I  
FLAGS($acD)  
$pc++
```

# LSR16



## Format:

LSR16     \$acR

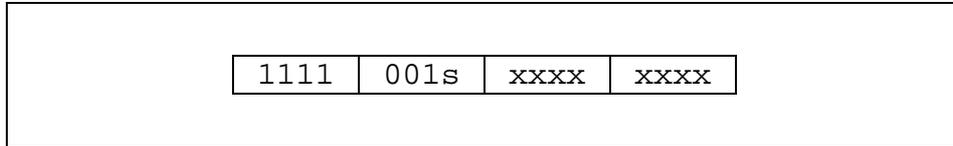
## Description:

Logically shifts right accumulator \$acR by 16.

## Operation:

```
$acR >>= 16  
FLAGS($acD)  
$pc++
```

# MADD



## Format:

MADD     \$axS.l, \$axS.h

## Description:

Multiply low part \$axS.l of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed) and add result to product register.

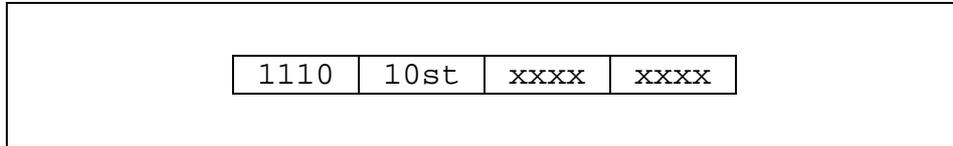
## Operation:

```
$prod += $axS.l * $axS.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MADDc



## Format:

MADDc    \$acS.m, \$axT.h

## Description:

Multiply middle part of accumulator \$acS.m by high part of secondary accumulator \$axT.h (treat them both as signed) and add result to product register.

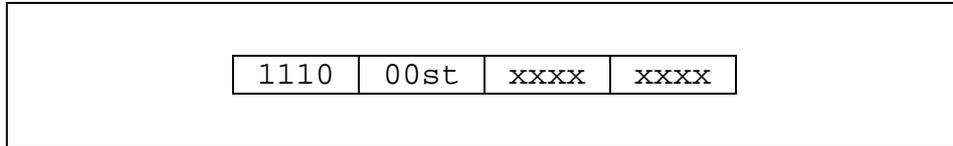
## Operation:

```
$prod += $acS.m * $axT.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MADDX



## Format:

MADDX      $\$(0x18+S*2), \$(0x19+T*2)$

## Description:

Multiply one part of secondary accumulator \$ax0 (selected by S) by one part of secondary accumulator \$ax1 (selected by T) (treat them both as signed) and add result to product register.

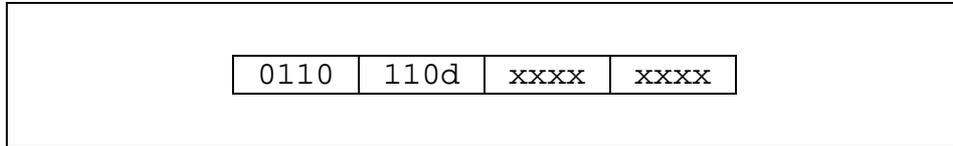
## Operation:

```
$prod += $(0x18+S*2) * $(0x19+T*2)
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MOV



## Format:

MOV            \$acD, \$ac(1-D)

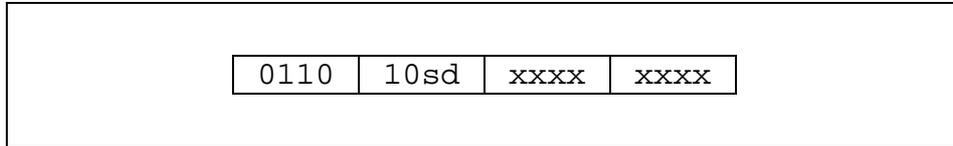
## Description:

Moves accumulator \$ax(1-D) to accumulator \$axD.

## Operation:

\$acD = \$ax(1-D)  
FLAGS(\$acD)  
\$pc++

# MOVAX



## Format:

MOVAX     \$acD, \$axS

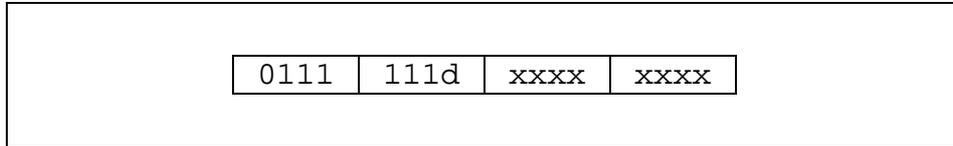
## Description:

Moves secondary accumulator \$axS to accumulator \$acD.

## Operation:

\$acD = \$axS  
FLAGS(\$acD)  
\$pc++

# MOVNP



## Format:

MOVNP     \$acD

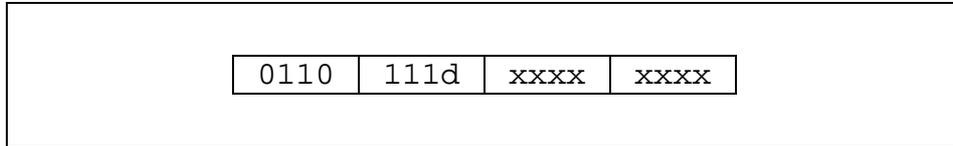
## Description:

Moves negative of multiply product from \$prod register to accumulator \$acD register.

## Operation:

```
$acD = -$prod  
FLAGS($acD)  
$pc++
```

# MOVP



## Format:

MOVP      \$acD

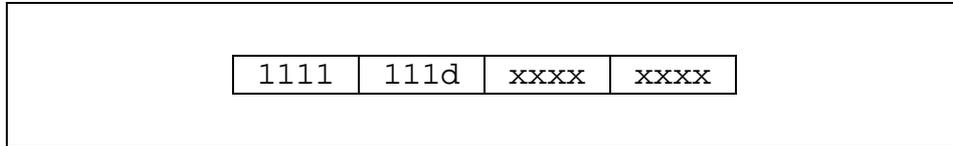
## Description:

Moves multiply product from \$prod register to accumulator \$acD register.

## Operation:

```
$acD = $prod  
FLAGS($acD)  
$pc++
```

# MOVPZ



## Format:

MOVPZ     \$acD

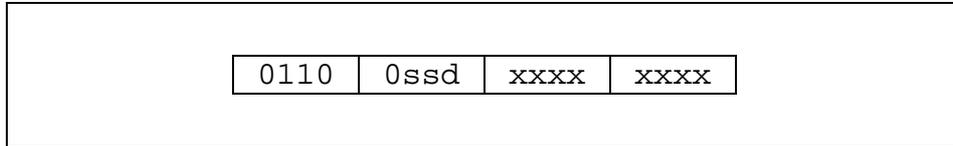
## Description:

Moves multiply product from \$prod register to accumulator \$acD register and sets \$acD.l to 0

## Operation:

```
$acD.hm = $prod.hm  
$acD.l = 0  
FLAGS($acD)  
$pc++
```

# MOVR



## Format:

MOVR                    \$acD, \$(0x18+S)

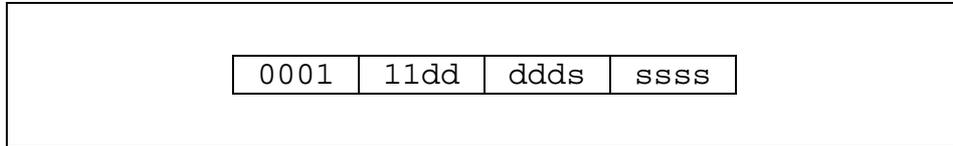
## Description:

Moves register \$(0x18+S) (sign extended) to middle accumulator \$acD.hm. Sets \$acD.l to 0.

## Operation:

```
$acD.hm = $(0x18+S)
$acD.l = 0
FLAGS($acD)
$pc++
```

# MRR



## Format:

MRR      \$D, \$S

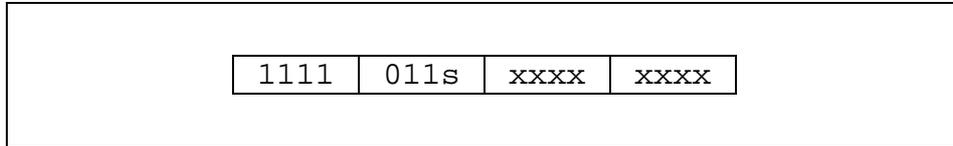
## Description:

Move value from register \$S to register \$D. Perform additional operation depending on destination register.

## Operation:

\$D = \$S  
\$pc++

# MSUB



## Format:

MSUB      \$axS.l, \$axS.h

## Description:

Multiply low part \$axS.l of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed) and subtract result from product register.

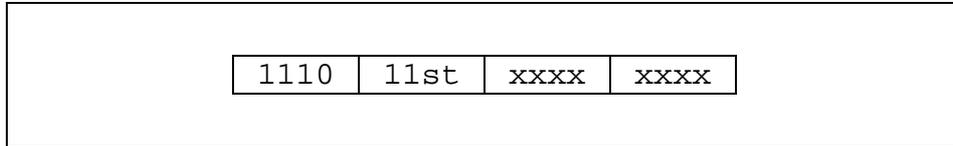
## Operation:

```
$prod -= $axS.l * $axS.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MSUBC



## Format:

MSUBC    \$acS.m, \$axT.h

## Description:

Multiply middle part of accumulator \$acS.m by high part of secondary accumulator \$axT.h (treat them both as signed) and subtract result from product register.

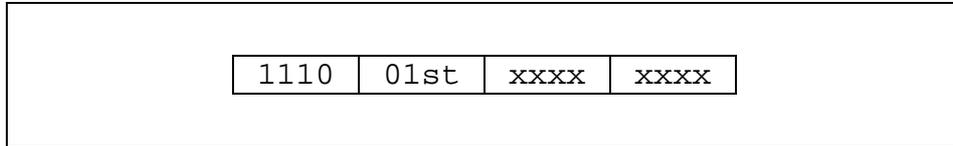
## Operation:

```
$prod -= $acS.m * $axT.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MSUBX



## Format:

MSUBX      $\$(0x18+S*2), \$(0x19+T*2)$

## Description:

Multiply one part of secondary accumulator \$ax0 (selected by S) by one part of secondary accumulator \$ax1 (selected by T) (treat them both as signed) and subtract result from product register.

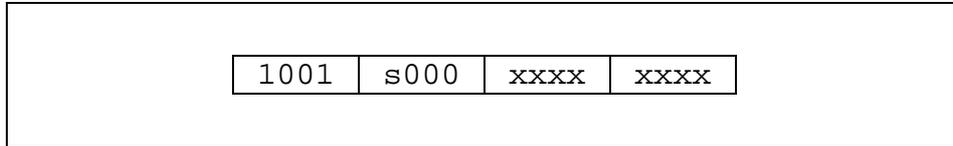
## Operation:

$\$prod -= \$(0x18+S*2) * \$(0x19+T*2)$   
\$pc++

## See also:

\$sr.AM bit affects multiply result

# MUL



## Format:

MUL      \$axS.l, \$axS.h

## Description:

Multiply low part \$axS.l of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

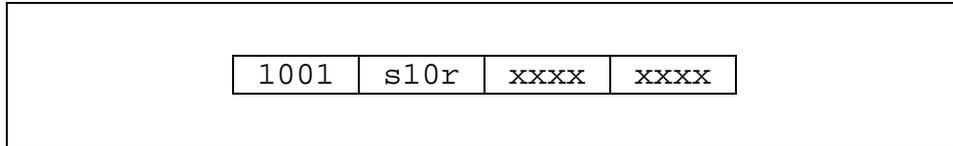
## Operation:

```
$prod = $axS.l * $axS.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULAC



## Format:

MULAC     \$axS.l, \$axS.h, \$acR

## Description:

Add product register to accumulator register \$acR. Multiply low part \$axS.l of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

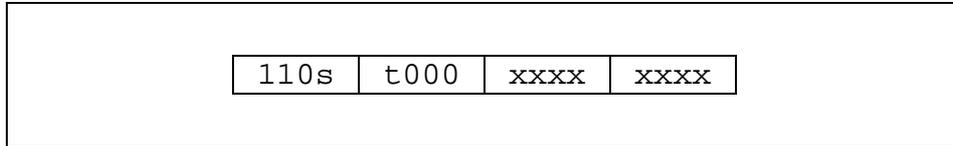
## Operation:

```
$acR += $prod
$prod = $axS.l * $axS.h
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULC



## Format:

MULC      \$acS.m, \$axT.h

## Description:

Multiply mid part of accumulator register \$acS.m by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

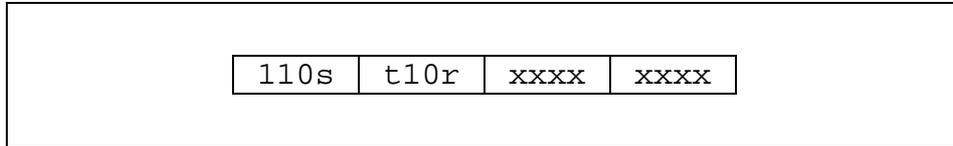
## Operation:

```
$prod = $acS.m * $axS.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULCAC



## Format:

MULCAC    \$acS.m, \$axT.h, \$acR

## Description:

Multiply mid part of accumulator register \$acS.m by high part \$axS.h of secondary accumulator \$axS (treat them both as signed). Add product register before multiplication to accumulator \$acR.

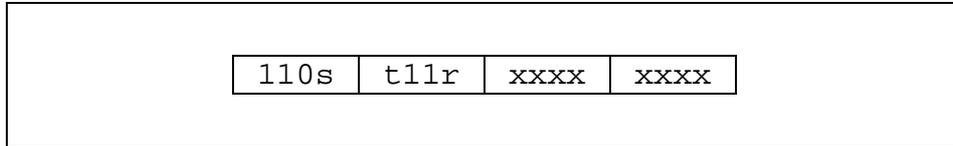
## Operation:

```
temp = $prod
$prod = $acS.m * $axS.h
$acR += temp
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULCMV



## Format:

MULCMV \$acS.m, \$axT.h, \$acR

## Description:

Multiply mid part of accumulator register \$acS.m by high part \$axS.h of secondary accumulator \$axS (treat them both as signed). Move product register before multiplication to accumulator \$acR.

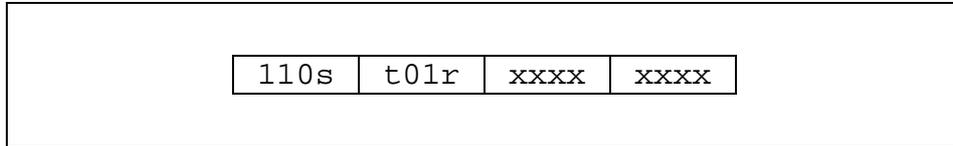
## Operation:

```
temp = $prod
$prod = $acS.m * $axS.h
$acR = temp
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULCMVZ



## Format:

MULCMVZ \$acS.m, \$axT.h, \$acR

## Description:

Multiply mid part of accumulator register \$acS.m by high part \$axS.h of secondary accumulator \$axS (treat them both as signed). Move product register before multiplication to accumulator \$acR, set low part of accumulator \$acR.l to zero.

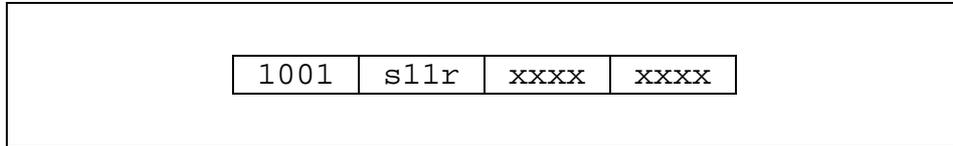
## Operation:

```
temp = $prod
$prod = $acS.m * $axS.h
$acR.hm = temp.hm
$acR.l = 0
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULMV



## Format:

MULMV    \$axS.l, \$axS.h, \$acR

## Description:

Move product register to accumulator register \$acR. Multiply low part \$axS.l of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

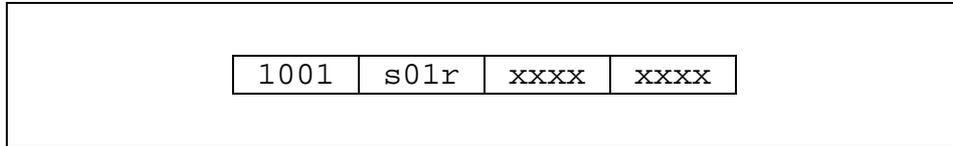
## Operation:

```
$acR = $prod  
$prod = $axS.l * $axS.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULMVZ



## Format:

MULMVZ    \$axS.l, \$axS.h, \$acR

## Description:

Move product register to accumulator register \$acR and clear low part of accumulator register \$acR.l. Multiply low part \$axS.l of secondary accumulator \$axS by high part \$axS.h of secondary accumulator \$axS (treat them both as signed).

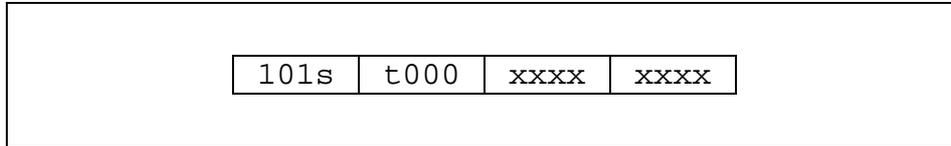
## Operation:

```
$acR.hm = $prod.hm  
$acR.l = 0  
$prod = $axS.l * $axS.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULX



## Format:

MULX      \$ax0.S, \$ax1.T

## Description:

Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

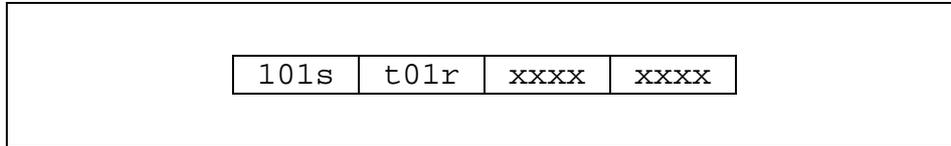
## Operation:

```
$prod = (S==0)?$ax0.l:ax0.h * (T==0)?$ax1.l:$ax1.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULXAC



## Format:

MULXAC \$ax0.S, \$ax1.T, \$acR

## Description:

Add product register to accumulator register \$acR. Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

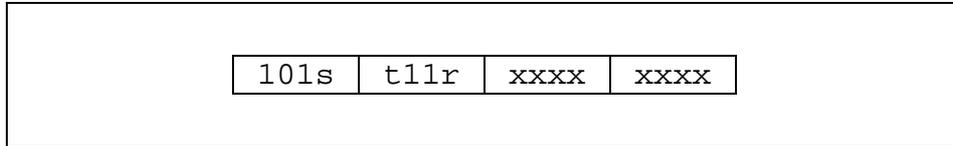
## Operation:

```
$acR += $prod  
$prod = (S==0)?$ax0.l:ax0.h * (T==0)?$ax1.l:$ax1.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULXMV



## Format:

MULXMV \$ax0.S, \$ax1.T, \$acR

## Description:

Move product register to accumulator register \$acR. Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

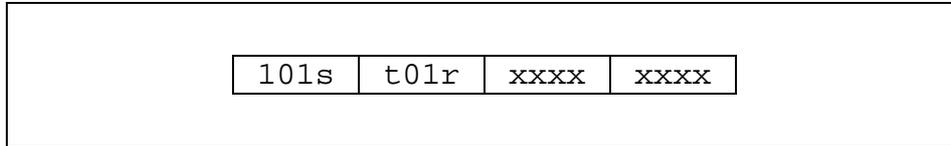
## Operation:

```
$acR = $prod  
$prod = (S==0)?$ax0.l:ax0.h * (T==0)?$ax1.l:$ax1.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# MULXMVZ



## Format:

MULXMV \$ax0.S, \$ax1.T, \$acR

## Description:

Move product register to accumulator register \$acR and clear low part of accumulator register \$acR.l. Multiply one part \$ax0 by one part \$ax1 (treat them both as signed). Part is selected by S and T bits. Zero selects low part, one selects high part.

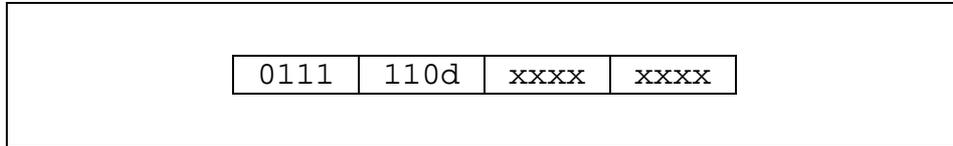
## Operation:

```
$acR.hm = $prod.hm  
$acR.l = 0  
$prod = (S==0)?$ax0.l:ax0.h * (T==0)?$ax1.l:$ax1.h  
$pc++
```

## See also:

\$sr.AM bit affects multiply result

# NEG



## Format:

NEG      \$acD

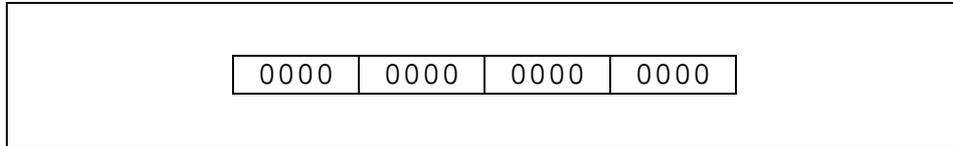
## Description:

Negate accumulator \$acD.

## Operation:

```
$acD =- $acD  
FLAGS($acD)  
$pc++
```

# NOP



## Format:

NOP

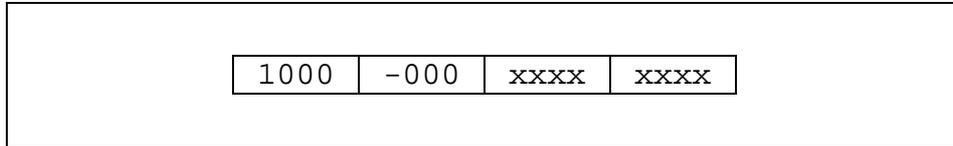
## Description:

No operation.

## Operation:

`$pc++;`

# NX



## Format:

NX

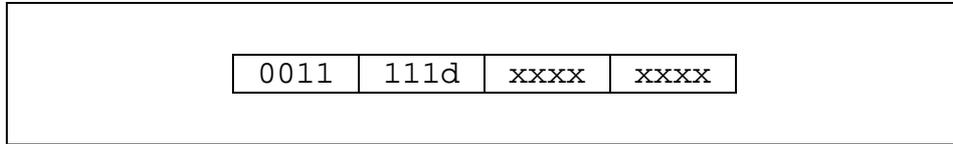
## Description:

No operation, but can be extended with extended opcode.

## Operation:

$\$pc++;$

# ORC



## Format:

ORC      \$acD.m, \$ac(1-D).m

## Description:

Logic OR middle part of accumulator \$acD.m with middle part of accumulator \$ax(1-D).m.

## Operation:

```
$acD.m |= $ac(1-D).m  
FLAGS($acD)  
$pc++
```

# ORI

0000	001r	0110	0000
iii	iii	iii	iii

## Format:

ORI            \$acD.m, #I

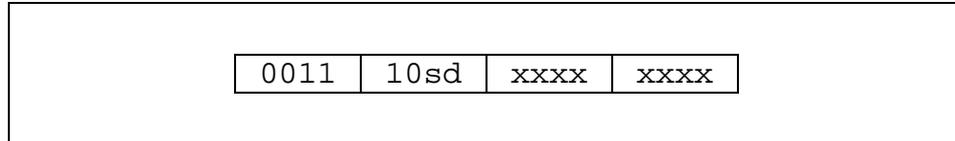
## Description:

Logic OR of accumulator mid part \$acD.m with immediate value I.

## Operation:

```
$acD.m |= #I  
FLAGS($acD)  
$pc++
```

# ORR



## Format:

ORR      \$acD.m, \$axS.h

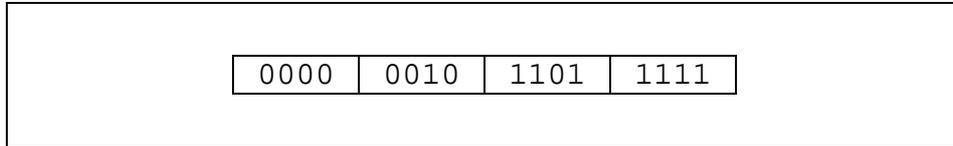
## Description:

Logic OR middle part of accumulator \$acD.m with high part of secondary accumulator \$axS.h.

## Operation:

```
$acD.m |= $axS.h  
FLAGS($acD)  
$pc++
```

# RET



## Format:

RET

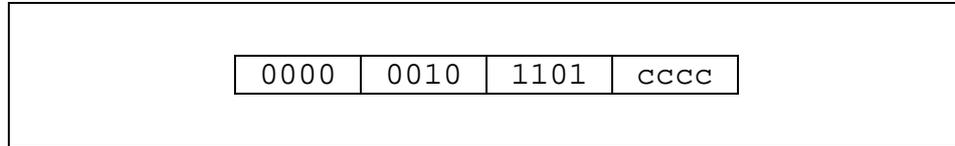
## Description:

Return from subroutine. Pops stored PC from call stack \$st0 and sets \$pc to this location.

## Operation:

```
$pc = $st0  
POP_STACK($st0)
```

# RETcc



## Format:

RETcc

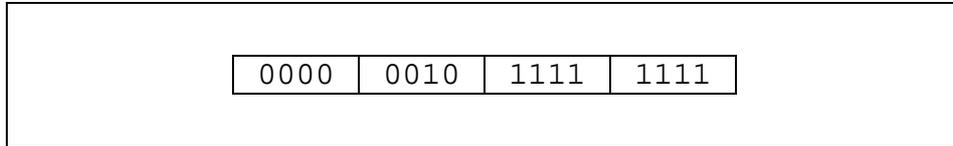
## Description:

Return from subroutine if condition cc has been met. Pops stored PC from call stack \$st0 and sets \$pc to this location.

## Operation:

```
IF (cc)    $pc = POP_STACK($st0)
ELSE      $pc += 2
```

# RTI



## Format:

RTI

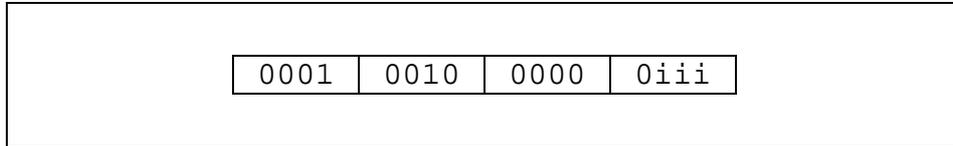
## Description:

Return from exception. Pops stored status register \$sr from data stack \$st1 and program counter PC from call stack \$st0 and sets \$pc to this location.

## Operation:

```
$sr = $st1  
POP_STACK($st1)  
$pc = $st0  
POP_STACK($st0)
```

# SBSET



## Format:

SBSET #I

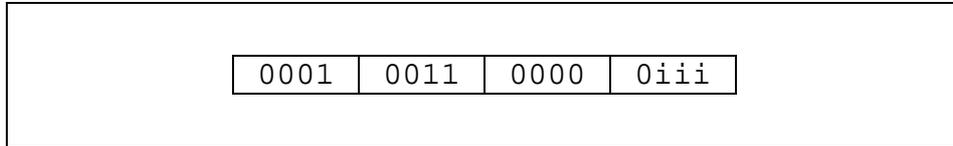
## Description:

Set bit of status register \$sr. Bit number is calculated by adding 6 to immediate value I.

## Operation:

```
$sr |= (I + 6)
$pc++
```

# SBCLR



## Format:

SBCLR #I

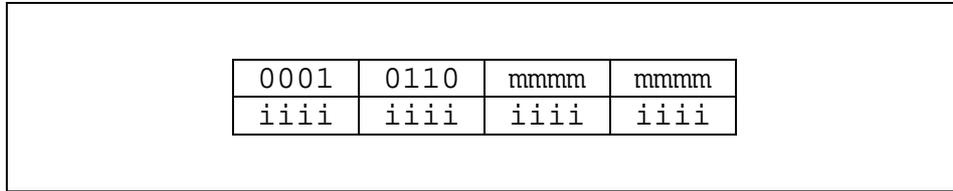
## Description:

Clear bit of status register \$sr. Bit number is calculated by adding 6 to immediate value I.

## Operation:

```
$sr &= ~(I + 6)  
$pc++
```

# SI



## Format:

SI            @M, #I

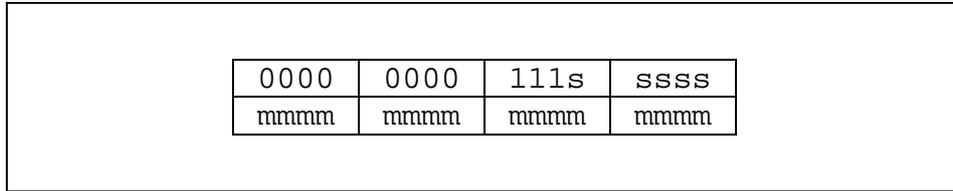
## Description:

Store 16-bit immediate value I to a memory location pointed by address M (M is 8-bit value sign extended).

## Operation:

MEM[M] = I  
\$pc += 2

# SR



## Format:

SR            @M, \$S

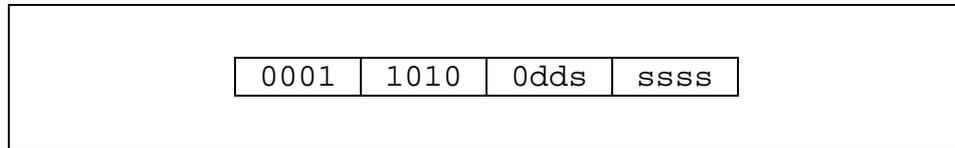
## Description:

Store value from register \$S to a memory pointed by address M.  
Perform additional operation depending on destination register.

## Operation:

MEM[M] = \$S  
\$pc += 2

# SRR



## Format:

SRR      @\$D, \$S

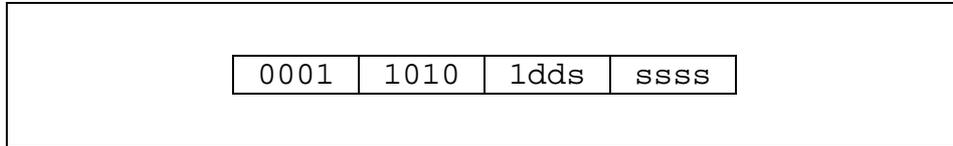
## Description:

Store value from source register \$S to a memory location pointed by addressing register \$D. Perform additional operation depending on source register.

## Operation:

MEM[\$D] = \$S  
\$pc++

# SRRD



## Format:

SRRD      @\$D, \$S

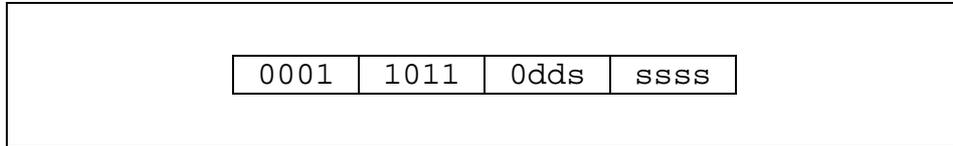
## Description:

Store value from source register \$S to a memory location pointed by addressing register \$D. Decrement register \$D. Perform additional operation depending on source register.

## Operation:

MEM[\$D] = \$S  
\$D--  
\$pc++

# SRRI



## Format:

SRRI      @\$D, \$S

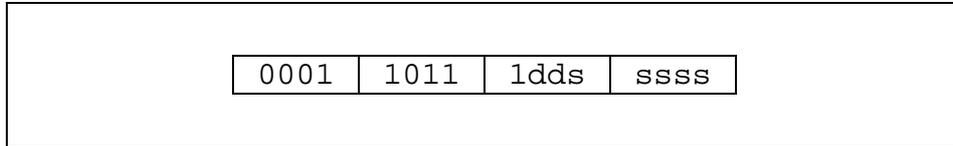
## Description:

Store value from source register \$S to a memory location pointed by addressing register \$D. Increment register \$D. Perform additional operation depending on source register.

## Operation:

MEM[\$D] = \$S  
\$D++  
\$pc++

# SRRN



## Format:

SRRN      @\$D, \$S

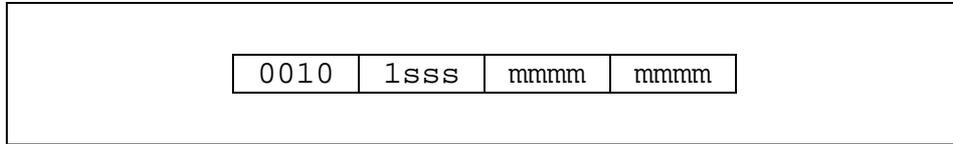
## Description:

Store value from source register \$S to a memory location pointed by addressing register \$D. Add indexing register \$(0x4+D) to register \$D. Perform additional operation depending on source register.

## Operation:

```
MEM[$D] = $S
$D += $(4+D)
$pc++
```

# SRS



## Format:

SRS @M, \$(0x18+S)

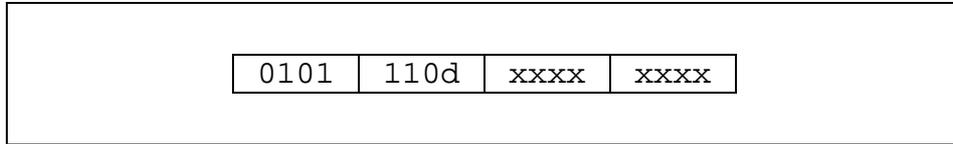
## Description:

Store value from register \$(0x18+S) to a memory pointed by address M. (8-bit sign extended). Perform additional operation depending on destination register.

## Operation:

MEM[M] = \$(0x18+S)  
\$pc += 2

# SUB



## Format:

SUB        \$acD, \$ac(1-D)

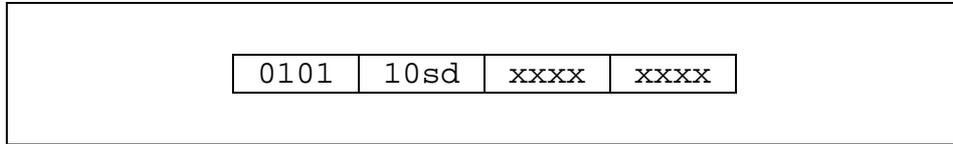
## Description:

Subtracts accumulator \$ac(1-D) from accumulator register \$acD.

## Operation:

```
$acD -= $ac(1-D)
FLAGS($acD)
$pc++
```

# SUBAX



## Format:

SUBAX     \$acD, \$axS

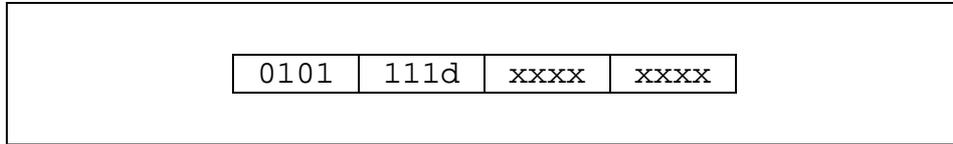
## Description:

Subtracts secondary accumulator \$axS from accumulator register \$acD.

## Operation:

```
$acD -= $axS  
FLAGS($acD)  
$pc++
```

# SUBP



## Format:

SUBP      \$acD

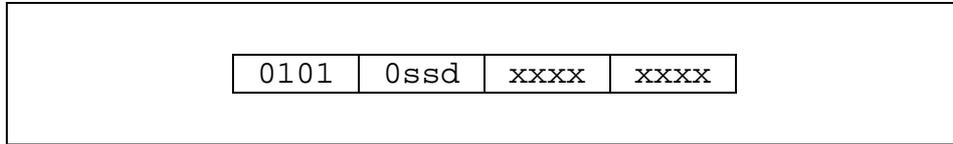
## Description:

Subtracts product register from accumulator register.

## Operation:

```
$acD -= $prod  
FLAGS($acD)  
$pc++
```

# SUBR



## Format:

SUBR      \$acD, \$(0x18+S)

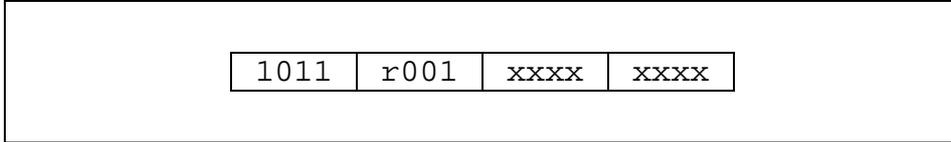
## Description:

Subtracts register \$(0x18+S) from accumulator \$acD register.

## Operation:

```
$acD -= $(0x18+S)
FLAGS($acD)
$pc++
```

# TST



## Format:

TST      \$acR

## Description:

Test accumulator \$acR

## Operation:

FLAGS (\$acR)  
\$pc++

# TSTAXH

1000	011r	xxxx	xxxx
------	------	------	------

## Format:

TST        \$axR.h

## Description:

Test high part of secondary accumulator \$axR.h.

## Operation:

FLAGS ( \$axR.h )  
\$pc++

# XORI

0000	001r	0010	0000
iii	iii	iii	iii

## Format:

XORI      \$acD.m, #I

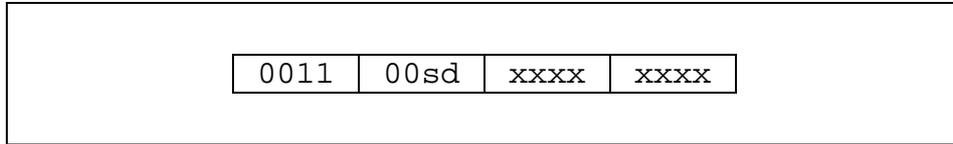
## Description:

Logic exclusive or (XOR) of accumulator mid part \$acD.m with immediate value I.

## Operation:

```
$acD.m ^= #I  
FLAGS($acD)  
$PC++
```

# XORR



## Format:

XORR      \$acD.m, \$axS.h

## Description:

Logic XOR (exclusive or) middle part of accumulator \$acD.m with high part of secondary accumulator \$axS.h.

## Operation:

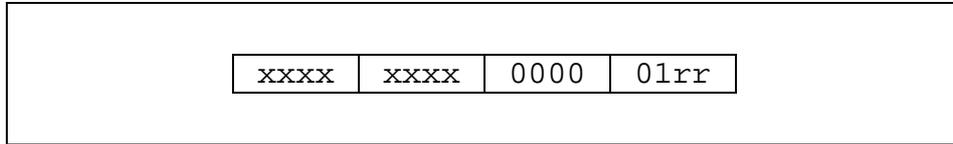
```
$acD.m ^= $axS.h  
FLAGS($acD)  
$pc++
```



## **6. *Extended opcodes decoding***

Extended opcodes do not exist on their own. These opcodes can only be attached to opcodes that allow extending (8 lower bits of opcode not used by opcode). Extended opcodes do not modify program counter \$pc register.

# 'DR



## Format:

'DR      \$arR

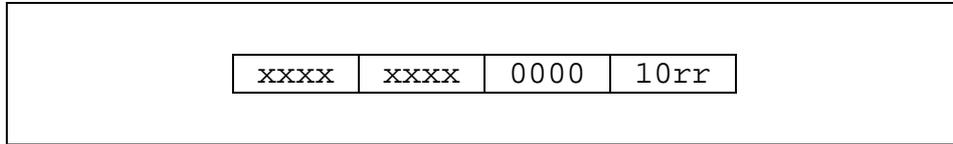
## Description:

Decrement addressing register \$arR.

## Operation:

\$arR ←

# 'IR



## Format:

'IR      \$arR

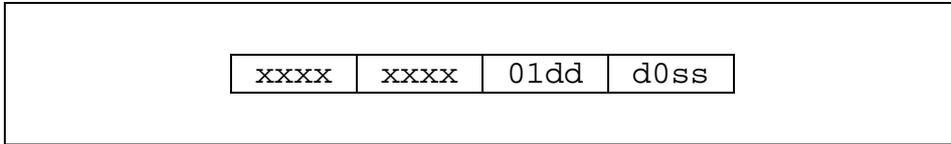
## Description:

Increment addressing register \$arR.

## Operation:

\$arR++

**'L**



**Format:**

'L           \$(0x18+D), @\$S

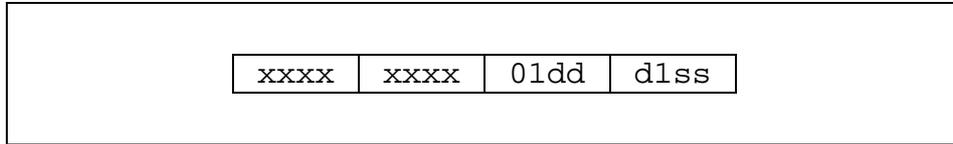
**Description:**

Load register \$(0x18+D) with value from memory pointed by register \$S. Post increment register \$S.

**Operation:**

\$(0x18+D) = MEM[\$S]  
\$S++

# 'LN



## Format:

'LN           \$(0x18+D), @\$S

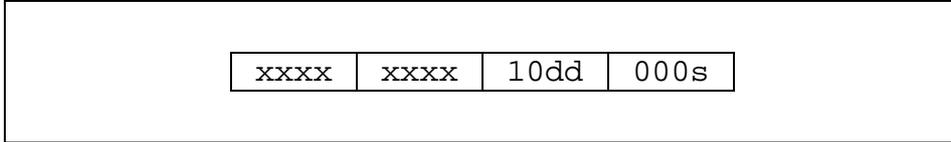
## Description:

Load register \$(0x18+D) with value from memory pointed by register \$S. Add indexing register register \$(0x04+S) to register \$S.

## Operation:

```
$(0x18+D) = MEM[$S]  
$S += $(0x04+S)
```

# 'LS



## Format:

'LS \$(0x18+D), \$acS.m

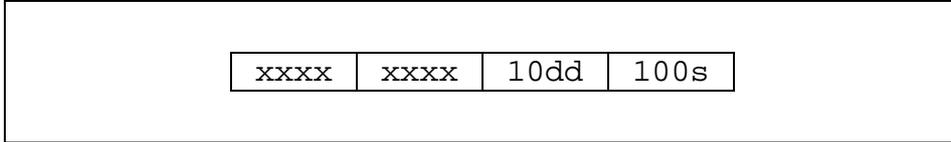
## Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Increment both \$ar0 and \$ar3.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0++  
$ar3++
```

# 'LSM



## Format:

'LSM      \$(0x18+D), \$acS.m

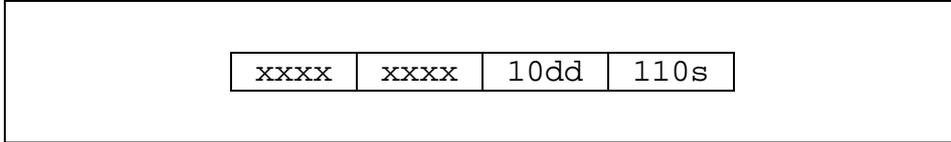
## Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Add corresponding indexing register \$ix3 to addressing register \$ar3 and increment \$ar0.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0++  
$ar3 += $ix3
```

# 'LSMN



## Format:

'LSMN      \$(0x18+D), \$acS.m

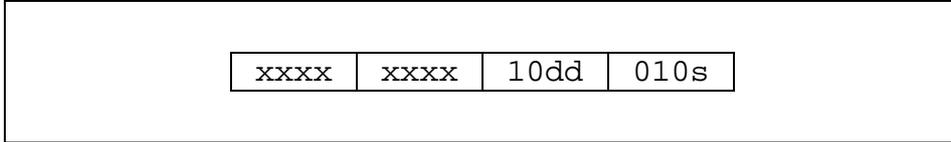
## Description:

Load register \$(0x18+D) with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and add corresponding indexing register \$ix3 to addressing register \$ar3.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0 += $ix0  
$ar3 += $ix3
```

# 'LSN



## Format:

'LSN       $\$(0x18+D)$ , \$acS.m

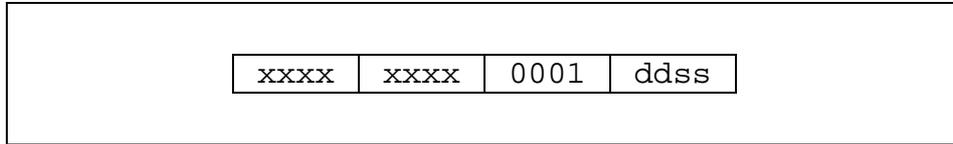
## Description:

Load register  $\$(0x18+D)$  with value from memory pointed by register \$ar0. Store value from register \$acS.m to memory location pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and increment \$ar3.

## Operation:

```
 $\$(0x18+D) = \text{MEM}[\$ar0]$   
 $\text{MEM}[\$ar3] = \$acS.m$   
 $\$ar0 += \$ix0$   
 $\$ar3++$ 
```

# 'MV



## Format:

'MV      \$(0x18+D), \$(0x1c+S)

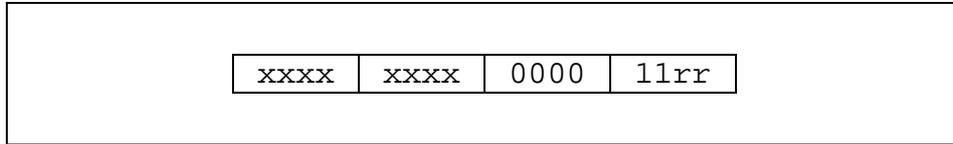
## Description:

Move value of register \$(0x1c+S) to the register \$(0x18+D).

## Operation:

$\$(0x18+D) = \$(0x1c+S)$

# 'NR



## Format:

'NR      \$arR

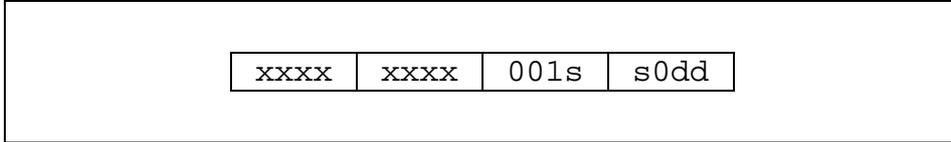
## Description:

Add corresponding indexing register \$ixR to addressing register \$arR.

## Operation:

\$arR += \$ixR

# 'S



## Format:

'S            @\$D, \$(0x1c+D)

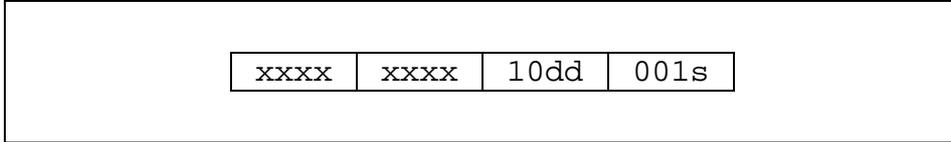
## Description:

Store value of register \$(0x1c+S) in the memory pointed by register \$D.  
Post increment register \$D.

## Operation:

MEM[\$D] = \$(0x1c+D)  
\$S++

# 'SL



## Format:

'SL            \$acS.m, \$(0x18+D)

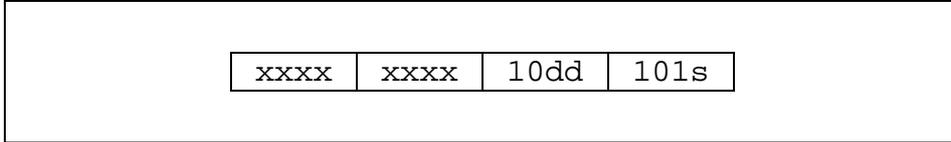
## Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Increment both \$ar0 and \$ar3.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0++  
$ar3++
```

# 'SLM



## Format:

'SLM      \$acS.m, \$(0x18+D)

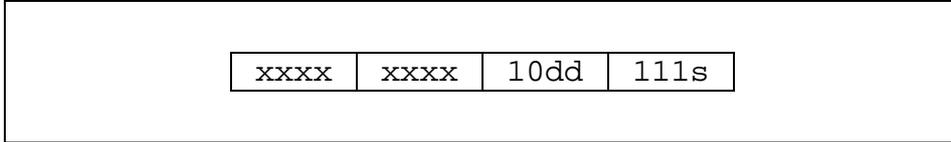
## Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix3 to addressing register \$ar3 and increment \$ar0.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0++  
$ar3 += $ix3
```

# 'SLMN



## Format:

'SLMN      \$acS.m, \$(0x18+D)

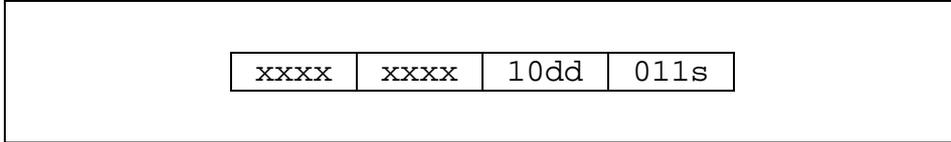
## Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and add corresponding indexing register \$ix3 to addressing register \$ar3.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0 += $ix0  
$ar3 += $ix3
```

# 'SLN



## Format:

'SLN      \$acS.m, \$(0x18+D)

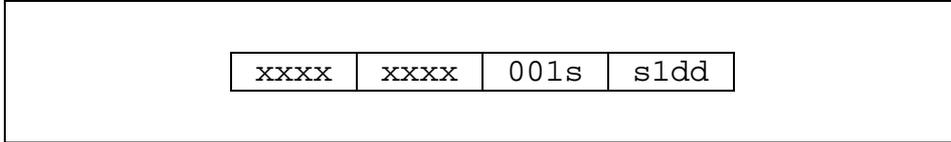
## Description:

Store value from register \$acS.m to memory location pointed by register \$ar0. Load register \$(0x18+D) with value from memory pointed by register \$ar3. Add corresponding indexing register \$ix0 to addressing register \$ar0 and increment \$ar3.

## Operation:

```
$(0x18+D) = MEM[$ar0]  
MEM[$ar3] = $acS.m  
$ar0 += $ix0  
$ar3++
```

# 'SN



## Format:

'SN            @\$D, \$(0x1c+D)

## Description:

Store value of register  $$(0x1c+S)$  in the memory pointed by register  $\$D$ .  
Add indexing register register  $$(0x04+D)$  to register  $\$D$ .

## Operation:

MEM[ $\$D$ ] =  $$(0x1c+D)$   
 $\$D$  +=  $$(0x04+D)$

## 7. Opcodes sorted by bit decoding

NOP		* 0000 0000 0000 0000
DAR		* 0000 0000 0000 01aa
IAR		* 0000 0000 0000 10aa
XXX	NOT USED	0000 0000 0000 11xx
ADDARN		* 0000 0000 0001 bbaa
HALT		* 0000 0000 0010 0001
LOOP		* 0000 0000 010r rrrr
BLOOP		* 0000 0000 011r rrrr
LRI		* 0000 0000 100r rrrr iiii iiii iiii iiii
XXX	NOT USED	0000 0000 101x xxxx
LR		* 0000 0000 110r rrrr mmmm mmmm mmmm mmmm
SR		* 0000 0000 111r rrrr mmmm mmmm mmmm mmmm
IF cc		* 0000 0010 0111 cccc
JMP cc		* 0000 0010 1001 cccc
CALL cc		* 0000 0010 1011 cccc
RET cc		* 0000 0010 1101 cccc
ADDI		* 0000 001r 0000 0000 iiii iiii iiii iiii
XORI		* 0000 001r 0010 0000 iiii iiii iiii iiii
ANDI		* 0000 001r 0100 0000 iiii iiii iiii iiii
ORI		* 0000 001r 0110 0000 iiii iiii iiii iiii
CMPI		* 0000 001r 1000 0000 iiii iiii iiii iiii
ANDCF		* 0000 001r 1010 0000 iiii iiii iiii iiii
ANDF		* 0000 001r 1100 0000 iiii iiii iiii iiii
ILRR		* 0000 001r 0001 mmaa
ADDIS		* 0000 010d iiii iiii
CMPIS		* 0000 011d iiii iiii
LRIS		* 0000 1rrr iiii iiii
LOOPI		* 0001 0000 iiii iiii aaaa aaaa aaaa aaaa
BLOOPI		* 0001 0001 iiii iiii aaaa aaaa aaaa aaaa
SBSET	bit set	* 0001 0010 ???? ?iii
SBCLR	bit clear	* 0001 0011 ???? ?iii
LSL/LSR		* 0001 010r 0sss ssss
ASL/ASR		* 0001 010r 1sss ssss
SI		* 0001 0110 iiii iiii mmmm mmmm mmmm mmmm
CALLR		* 0001 0111 rrr1 1111
JMPR		* 0001 0111 rrr0 1111

LRR(I D X)		* 0001 100x xaar rrrr
SRR(I D X)		* 0001 101x xaar rrrr
MRR		* 0001 11dd ddds ssss
LRS		* 0010 0rrr mmmm mmmm
SRS		* 0010 1rrr mmmm mmmm
XORR		* 0011 00sr xxxx xxxx
ANDR		* 0011 01sr xxxx xxxx
ORR		* 0011 10sr xxxx xxxx
ANDC		* 0011 110r xxxx xxxx
ORC		* 0011 111r xxxx xxxx
ADDR		* 0100 0ssd xxxx xxxx
ADDAX		* 0100 10sd xxxx xxxx
ADD		* 0100 110d xxxx xxxx
ADDP		* 0100 111d xxxx xxxx
SUBR		* 0101 0ssd xxxx xxxx
SUBAX		* 0101 10sd xxxx xxxx
SUB		* 0101 110d xxxx xxxx
SUBP		* 0101 111d xxxx xxxx
MOVR		* 0110 0ssd xxxx xxxx
MOVAX		* 0110 10sd xxxx xxxx
MOV		* 0110 110d xxxx xxxx
MOV P		* 0110 111d xxxx xxxx
ADDAXL		* 0111 00sr xxxx xxxx
INCM		* 0111 010r xxxx xxxx
INC		* 0111 011r xxxx xxxx
DECM		* 0111 100r xxxx xxxx
DEC		* 0111 101r xxxx xxxx
NEG		* 0111 110r xxxx xxxx
MOVNP		* 0111 111r xxxx xxxx
NX		1000 x000 xxxx xxxx
CLR		* 1000 x001 xxxx xxxx
CMP		* 1000 0010 xxxx xxxx
???	UNUSED	1000 0011 xxxx xxxx
CLRP		* 1000 0100 xxxx xxxx
TSTAXH		* 1000 011x xxxx xxxx
M0/M2		1000 101x xxxx xxxx
CLR15/SET15		1000 110x xxxx xxxx
SET40/16		1000 111x xxxx xxxx
MUL		* 1001 a000 xxxx xxxx

ASR16	* 1001 r001 xxxx xxxx
MULMVZ	* 1001 a01r xxxx xxxx
MULAC	* 1001 a10r xxxx xxxx
MULMV	* 1001 a11r xxxx xxxx
MULX	* 101b a000 xxxx xxxx
???	1010 r001 xxxx xxxx
TST	* 1011 r001 xxxx xxxx
MULXMVZ	* 101b a01r xxxx xxxx
MULXAC	* 101b a10r xxxx xxxx
MULXMV	* 101b a11r xxxx xxxx
MULC	* 110s a000 xxxx xxxx
CMP	* 110x r001 xxxx xxxx
MULCMVZ	* 110s a01r xxxx xxxx
MULCAC	* 110s a10r xxxx xxxx
MULCMV	* 110s a11r xxxx xxxx
MADDX	** 1110 00st xxxx xxxx
MSUBX	** 1110 01st xxxx xxxx
MADDC	** 1110 10st xxxx xxxx
MSUBC	** 1110 11st xxxx xxxx
LSL16	* 1111 000r xxxx xxxx
MADD	* 1111 001s xxxx xxxx
LSR16	* 1111 010r xxxx xxxx
MSUB	* 1111 011s xxxx xxxx
ADDPAXZ	* 1111 10ar xxxx xxxx
CLRL	* 1111 110r xxxx xxxx
MOVPZ	* 1111 111r xxxx xxxx

### Opcode Extensions

[D I N]R	* xxxx xxxx 0000 nnaa
MV	* xxxx xxxx 0001 ddss
S[N]	* xxxx xxxx 001r rnaa
L[N]	* xxxx xxxx 01dd diss
LS[NM M N]	* xxxx xxxx 10dd ba0r
SL[NM M N]	* xxxx xxxx 10dd balr
LD[NM M N]	xxxx xxxx 11mn barr
LD2[NM M N]	xxxx xxxx 11rm ball

## IX. References

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